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Variable elimination in binary CSPs

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Abstract

We investigate rules which allow variable elimination in binary CSP (constraint satisfaction problem) instances while conserving satisfiability. We study variable-elimination rules based on the language of forbidden patterns enriched with counting and quantification over variables and values. We propose new rules and compare them, both theoretically and experimentally. We give optimised algorithms to apply these rules and show that each define a novel tractable class. Using our variable-elimination rules in preprocessing allowed us to solve more benchmark problems than without.

1. Introduction

Constraint satisfaction provides a generic model for many NP-hard problems encountered in fields such as artificial intelligence, bioinformatics and operations research. In this paper, we study binary CSP instances, in which each constraint concerns at most two variables. It is well known that all CSP instances can be expressed as binary instances, via the dual encoding (Dechter & Pearl, 1989) or the hidden variable encoding (Rossi, Petrie, & Dhar, 1990).

Since the binary CSP is NP-complete, it is of practical interest to find polynomial-time operations which reduce the size of the search space. One obvious way to reduce search space size is by variable elimination.

Variable elimination is classic in those families of constraint problems in which variables can be eliminated without changing the nature of the constraints: we can cite Gaussian elimination in systems of linear equations over a field (Schrijver, 1999) or variable-elimination resolution in boolean formulae in CNF (Subbarayan & Pradhan, 2004). Indeed, any variable x_i can be eliminated from a *general-arity* CSP instance by joining all constraints whose scope includes x_i and projecting the resulting relation R with scope Y onto the variables $Y \setminus \{x_i\}$ (Dechter, 1999; Larrosa & Dechter, 2003). Call this relation R^{-x_i} . Unfortunately, this often introduces a high-arity constraint and this can be counterproductive in terms of both memory and time. Under certain conditions, a binary CSP instance will remain binary after this join-and-project variable elimination of x_i . For example, this is clearly the case if x_i is constrained by only two other variables since, in this case, R^{-x_i} is binary. A more interesting case is when all constraints with x_i in their scope share a majority polymorphism since, in this case, the relation R^{-x_i} is equivalent to the join of its binary projections (Jeavons, Cohen, & Cooper, 1998). Fourier's algorithm for variable elimination applied to a system of binary

linear inequalities (Koubarakis, 2006; Schrijver, 1999) can be viewed as just one example of this general rule, since binary linear inequalities are all closed under the majority polymorphism *median*. Another interesting case is when there is a functional constraint of the form $x_i = f(x_j)$ (where f is a function) for some other variable x_j : the relation R^{-x_i} is then equivalent to the join of its projections onto the pairs of variables (x_j, x_k) ($k \neq i, j$) (Zhang & Yap, 2011).

Unfortunately, the introduction of a large number of new constraints, even if they are still binary, may again be counterproductive. Therefore, we concentrate in this paper on rules which do not introduce new constraints when a variable is eliminated.

Various rules have been found which allow the elimination of a variable without introducing new constraints and without changing the satisfiability of the instance (Cohen, Cooper, Escamocher, & Zivny, 2015; Cooper, 2014; Cooper, Jeavons, & Salamon, 2010). Such rules were used, for example, in the deep optimisation solution to the spectrum repacking problem (Newman, Fréchette, & Leyton-Brown, 2018). Discovery of new variable-elimination rules may have not only practical but also theoretical applications. For example, simple rules for variable or value elimination are used by Beigel and Eppstein (Beigel & Eppstein, 1995) in their algorithms with low worst-case time bounds for such NP-complete problems as 3-COLOURING and 3SAT: these simplification operations are an essential first step before the use of decompositions into subproblems with smaller domains. In the theory of fixed-parameter tractability, variable elimination is often an essential ingredient of polynomial kernelisation algorithms. For example, in the Point Line Cover problem (find k straight lines which cover n points), if at least $k + 1$ points lie on a line, then they can be effectively eliminated since they must be covered by this line (Kratsch, Philip, & Ray, 2016). A form of variable elimination may also occur during the modelling phase. For example, in the modelling as a CSP of the determination of the structure of a molecule from its chemical formula and other information obtained from nuclear magnetic resonance spectroscopy, the position of hydrogen atoms are not modelled since their positions are uniquely determined by the multigraph of connections between the other atoms (Omriani & Naanaa, 2016).

We now define the notions that we will need in the rest of the paper.

Definition 1 A binary CSP instance $I = \langle X, \mathcal{D}, R \rangle$ comprises

- a set X of n variables x_1, \dots, x_n ,
- a domain $\mathcal{D}(x_i)$ for each variable x_i ($i = 1, \dots, n$), and
- a binary constraint relation R_{ij} for each pair of distinct variables x_i, x_j ($i, j \in \{1, \dots, n\}$).

For notational convenience, we assume that there is exactly one binary relation R_{ij} for each pair of variables. Thus, in the absence of an explicit constraint between x_i and x_j , we define R_{ij} to be $\mathcal{D}(x_i) \times \mathcal{D}(x_j)$. Furthermore, R_{ji} (viewed as a boolean matrix) is always the transpose of R_{ij} . We say that x_i *constrains* x_j if R_{ij} is different from $\mathcal{D}(x_i) \times \mathcal{D}(x_j)$, and we use e to denote the number of pairs of variables $\{x_i, x_j\}$ such that x_i constrains x_j . An assignment $\langle v_1, \dots, v_m \rangle$ to variables $\langle x_{i_1}, \dots, x_{i_m} \rangle$ is *consistent* if $v_j \in \mathcal{D}(x_{i_j})$ (for $j = 1, \dots, m$) and $(v_j, v_k) \in R_{i_j, i_k}$ (for all j, k such that $1 \leq j < k \leq m$). A *solution* to I is a consistent assignment to all variables in X . For notational convenience we can also view a solution as a mapping s from X to the union of the variable domains such that $\langle s(x_1), \dots, s(x_n) \rangle$ is consistent.

It will sometimes be convenient to associate a binary CSP instance with its microstructure, a labelled graph whose vertices are the variable-value assignments and which has positive and negative

edges. If $(v_i, v_j) \in R_{ij}$, we say that the assignments $\langle x_i, v_i \rangle, \langle x_j, v_j \rangle$ (or more simply v_i, v_j) are *compatible* and that $v_i v_j$ is a *positive edge*, otherwise v_i, v_j are *incompatible* and $v_i v_j$ is a *negative edge*. For simplicity of notation we can assume that variable domains are disjoint, so that using v_i as a shorthand for $\langle x_i, v_i \rangle$ is unambiguous. We say that $v_i \in \mathcal{D}(x_i)$ has a *support* at variable x_j if there exists $v_j \in \mathcal{D}(x_j)$ such that $v_i v_j$ is a positive edge. A binary CSP instance I is *arc consistent* if for all pairs of distinct variables x_i, x_j , each $v_i \in \mathcal{D}(x_i)$ has a support at x_j . Arc consistency is ubiquitous in constraint solvers: it is applied both before and during search in binary CSPs since it can be established in $O(ed^2)$ time, where e is the number of binary constraints and d the maximum domain size (Bessière, Régin, Yap, & Zhang, 2005).

In Section 2 we introduce formally the notion of variable-elimination rule in binary CSPs and give a known example (the \exists snake property) which we will compare theoretically and experimentally with novel variable elimination rules defined in this paper. Then, in Section 3 we define a stronger rule, called DE-snake, which subsumes the \exists snake rule. In Section 4 we give the definition of a simple variable elimination rule based on a triangle of variable-value assignments. It is well known that the broken-triangle property (Cooper et al., 2010) is a variable-elimination rule. In Section 5 we generalise broken triangles to broken polyhedra. In Section 6 we give a family of variable-elimination rules based on broken polyhedra of dimension k . On a more practical level, in Section 7 we define a variable-elimination rule, based on the absence of broken tetrahedra, which can be applied with the same worst-case time complexity as the broken-triangle rule but is strictly stronger. In Section 8 we show that most of the rules we have presented in this paper are theoretically incomparable. In Section 9 we present the results of our experimental trials on 3,557 benchmark instances. The variable-elimination rules allowed us to solve more instances when they were applied in a preprocessing step, but we recommend more research to better target those variables that are likely to be eliminated before integrating these rules in a general-purpose solver. In Section 10 we show that each of the variable-elimination rules presented in this paper allows us to define a tractable class which can be recognised in polynomial-time.

2. Variable-elimination rules

We study conditions under which a variable x_i can be eliminated from a binary CSP instance while conserving satisfiability. A simple example of such a condition is that there exists a value $v_i \in \mathcal{D}(x_i)$ which is compatible with all assignments to all other variables. Clearly any solution s to the instance I' obtained by eliminating x_i can be extended to a solution to the original instance I by setting $s(x_i) = v_i$. Another simple example is that the variable x_i has a singleton domain $\{v_i\}$. This second example demonstrates that when eliminating the variable x_i we need to retain the projections onto $X \setminus \{x_i\}$ of all constraints whose scope includes x_i , since in this example we must first eliminate from all domains $\mathcal{D}(x_j)$ ($j \neq i$) those values that are not compatible with $\langle x_i, v_i \rangle$. Thus, the instance I' obtained by *eliminating a variable* x_i from a binary CSP instance I is identical to I except that (1) $\forall j \neq i$, we have deleted from $\mathcal{D}(x_j)$ all values v_j such that $\langle x_j, v_j \rangle$ has no support at x_i in I , and (2) we have deleted the variable x_i and all constraints with x_i in their scope.

We require the following formal definition in order to study provably-correct variable-elimination rules (Cohen et al., 2015).

Definition 2 A satisfiability-conserving variable-elimination condition (or a var-elim condition) is a polytime-computable property $P(x_i)$ of a variable x_i in a binary CSP instance I such that when $P(x_i)$ holds the instance I' obtained from I by eliminating x_i is satisfiable if and only if I is

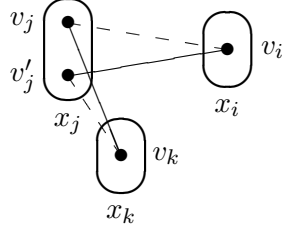


Figure 1: The \exists snake property says that this snake pattern does not occur on value v_i for variable x_i (for any variables x_j, x_k and any values v_j, v'_j, v_k).

satisfiable. Such a property $P(x_i)$ is a solution-conserving variable-elimination condition (sol-var-elim condition) if it is possible to construct a solution to I from any solution s' to I' in polynomial time.

A sol-var-elim condition not only allows us to eliminate variables while conserving satisfiability but also allows the polynomial-time recovery of at least one solution to the original instance I from a solution to the reduced instance I' . All the var-elim properties given in this paper are also sol-var-elim properties.

We end this section by giving an example of a known variable-elimination rule called the \exists snake rule. It is based on forbidding a pattern of positive and negative edges (shown in Figure 1) on one value v_i for x_i , the variable to be eliminated. In figures, broken lines represent negative edges (incompatible pairs) and solid lines represent positive edges (compatible pairs). A pattern is a (generally small) binary CSP instance in which the compatibility of certain values (such as v_i, v_k in Figure 1) may be left unspecified. A pattern P occurs in a CSP instance I if there is a homomorphism from P to I respecting variables and mapping positive edges to positive edges and negative edges to negative edges (Cohen et al., 2015). The \exists snake rule is one of the four variable-elimination rules based on forbidding an irreducible pattern on the variable to be eliminated (Cohen et al., 2015). Out of these four rules, we chose \exists snake to compare with the new rules presented in this paper since, among these four rules, it would appear to be the most promising in terms of time complexity and eliminating power.

Definition 3 *A variable x_i satisfies the \exists snake property if $\exists v_i \in \mathcal{D}(x_i)$ such that $\forall x_j \in X \setminus \{x_i\}$, $\forall v_j, v'_j \in \mathcal{D}(x_j)$, $\forall x_k \in X \setminus \{x_i, x_j\}$ $\forall v_k \in \mathcal{D}(x_k)$, we do not have $(v_i, v_j) \notin R_{ij}$, $(v_i, v'_j) \in R_{ij}$, $(v_j, v_k) \in R_{jk}$ and $(v'_j, v_k) \notin R_{jk}$.*

The proof of the following proposition can be found in Appendix A where we give an optimised algorithm, making use of appropriate data structures, to apply this variable-elimination rule until convergence.

Proposition 1 *Variable eliminations by the \exists snake property can be applied until convergence in $O(ed^3)$ time and $O(ed^2)$ space.*

As we will show in the rest of this paper, other variable-elimination rules can be found by enriching the language of forbidden patterns by allowing arbitrary quantification and counting. Previous work only considered quantification on values for the variable to be eliminated (Cohen et al., 2015).

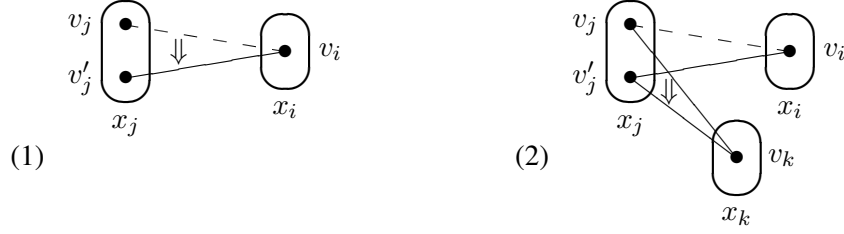


Figure 2: The DE-snake property says that for some value $v_i \in \mathcal{D}(x_i)$, for each value v_j incompatible with v_i , there is a value v'_j such that (1) v'_j is compatible with v_i , and (2) v'_j is compatible with all assignments v_k to a third variable x_k which are compatible with v_j .

3. Variable elimination by the DE-snake rule

We show in this section that the \exists snake rule is subsumed by a stronger rule that we call the DE-snake (double-existential snake) rule. It is again based on forbidding the snake pattern shown in Figure 1 but, compared to the \exists snake rule, has an existential (rather than universal) quantifier on the value v'_j .

Definition 4 A variable x_i satisfies the DE-snake property if $\exists v_i \in \mathcal{D}(x_i)$ such that $\forall x_j \in X \setminus \{x_i\}$, $\forall v_j \in \mathcal{D}(x_j)$ with $(v_i, v_j) \notin R_{ij}$, $\exists v'_j \in \mathcal{D}(x_j)$ such that (1) $(v'_j, v_i) \in R_{ji}$ and (2) $\forall x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$, we do not have $(v_j, v_k) \in R_{jk}$ and $(v'_j, v_k) \notin R_{jk}$.

The DE-snake property is illustrated in Figure 2. The intuition behind this property is that any solution to the instance obtained after elimination of x_i can be extended to a solution to the original instance by assigning v_i to x_i and changing those values v_j which are incompatible with v_i to some other value v'_j .

Theorem 1 The DE-snake property is a sol-var-elim condition in binary CSP instances.

Proof: Let $I = \langle X, \mathcal{D}, R \rangle$ be a binary CSP instance satisfying the DE-snake property on x_i . Let I' be the instance obtained by eliminating variable x_i from I . If there is no solution to I' , then obviously there is none to I . Now suppose that there is a solution s' to I' . We will show that I has a solution s . Our proof is constructive and there is an obvious polynomial-time algorithm to produce s from s' . Since x_i satisfies the DE-snake property, $\exists v_i \in \mathcal{D}(x_i)$ such that $\forall x_j \in X \setminus \{x_i\}$, $\forall v_j \in \mathcal{D}(x_j)$ with $(v_i, v_j) \notin R_{ij}$, $\exists u_j(v_j) \in \mathcal{D}(x_j)$ (i.e. there exists a value u_j which is a function of v_j) such that (1) $(u_j(v_j), v_i) \in R_{ji}$ and (2) $\forall x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$, we do not have $(v_j, v_k) \in R_{jk}$ and $(u_j(v_j), v_k) \notin R_{jk}$.

Let Y be the set of variables $x_j \in X \setminus \{x_i\}$ such that $(s'(x_j), v_i) \in R_{ji}$ and \bar{Y} the set of variables $x_j \in X \setminus \{x_i\}$ such that $(s'(x_j), v_i) \notin R_{ji}$. For each $x_j \in Y$, set $s(x_j) := s'(x_j)$. For each $x_j \in \bar{Y}$, set $s(x_j) := u_j(s'(x_j))$. Finally, set $s(x_i) := v_i$. By definition of s and $u_j(s'(x_j))$, we have $(s(x_j), s(x_i)) \in R_{ji}$ for each $x_j \in X \setminus \{x_i\}$. For pairs of variables $x_j, x_k \in X \setminus \{x_i\}$, we need to consider three cases:

1. If $x_j, x_k \in Y$, then clearly $(s(x_j), s(x_k)) \in R_{jk}$ since s' was a solution to I' .

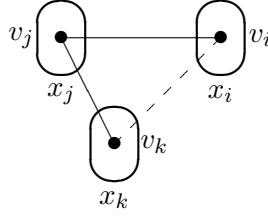


Figure 3: The open-triangle pattern. The variable x_i can be eliminated by the triangle property if there is some variable $x_j \neq x_i$ such that for all $v_j \in \mathcal{D}(x_j)$, there exists $v_i \in \mathcal{D}(x_i)$ such that $(v_j, v_i) \in R_{ji}$ and this open-triangle pattern does not occur on (v_i, v_j, v_k) for any v_k .

2. If $x_j \in \bar{Y}$ and $x_k \in Y$, then setting $v_j = s'(x_j)$ and $v_k = s'(x_k)$ in the definition of the DE-snake property, from condition (2) in this definition, we must have $(u_j(s'(x_j)), s'(x_k)) \in R_{jk}$ since $(s'(x_j), s'(x_k)) \in R_{jk}$. Hence $(s(x_j), s(x_k)) \in R_{jk}$ by definition of s .
3. If $x_j, x_k \in \bar{Y}$, then for exactly the same reason as in case 2, we again must have $(u_j(s'(x_j)), s'(x_k)) \in R_{jk}$. In other words, by definition of s , $(s(x_j), s'(x_k)) \in R_{jk}$. Now we apply again the definition of the DE-snake property but this time with the roles of the variables x_j, x_k reversed and with $v_j = s'(x_k)$ and $v_k = s(x_j)$: we can deduce that we must have $(u_k(s'(x_k)), s(x_j)) \in R_{kj}$. Thus, by definition of s , $(s(x_k), s(x_j)) \in R_{kj}$.

We have just shown that for all pairs of variables of I , s satisfies the binary constraint on this pair of variables. Hence s is a solution to I . ■

The following proposition shows that the worst-case complexity of applying the DE-snake rule is no worse than the complexity of applying the \exists snake rule given in Proposition 1. Its proof can be found in Appendix B where we give an optimised algorithm, making use of appropriate data structures, to apply the DE-snake variable-elimination rule until convergence.

Proposition 2 *Variable eliminations by the DE-snake property can be applied until convergence in $O(ed^3)$ time and $O(ed^2)$ space.*

4. Variable elimination by the triangle property

The variable-elimination rule presented in this section says that x_i can be eliminated if for some variable $x_j \neq x_i$, for all assignments $v_j \in \mathcal{D}(x_j)$ to y , in $I[\langle x_j, v_j \rangle]$ (the reduced instance consisting of the set of variable-value assignments compatible with $\langle x_j, v_j \rangle$) there is an assignment $\langle x_i, v_i \rangle$ compatible with all assignments to all variables $x_k \in X \setminus \{x_i, x_j\}$. In other words, there is some variable $x_j \neq x_i$ such that for all $v_j \in \mathcal{D}(x_j)$, there exists $v_i \in \mathcal{D}(x_i)$ such that $(v_j, v_i) \in R_{ji}$ and the open-triangle pattern shown in Figure 3 does not occur.

Definition 5 *A variable x_i satisfies the triangle property if $\exists x_j \in X \setminus \{x_i\}$ such that $\forall v_j \in \mathcal{D}(x_j)$, $\exists v_i \in \mathcal{D}(x_i)$ with $(v_i, v_j) \in R_{ij}$ such that $\forall x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$, $(v_j, v_k) \in R_{jk}$ implies that $(v_i, v_k) \in R_{ik}$.*

Although a variable satisfying the triangle property was originally known as ‘not Triangle-supported’ (Cooper, 2014), we use the name ‘triangle property’ in this paper for simplicity of presentation.

Theorem 2 *The triangle property is a sol-var-elim condition in binary CSP instances.*

Proof: Let I be a binary CSP instance satisfying the triangle property on x_i . Let s' be a solution to I' , the instance obtained by eliminating variable x_i from I . We will show that I has a solution s . Our proof is constructive and there is an obvious polynomial-time algorithm to produce s from s' . Since x_i satisfies the triangle property, $\exists x_j \in X \setminus \{x_i\}$ such that $\forall v_j \in \mathcal{D}(x_j), \exists v_i(v_j) \in \mathcal{D}(x_i)$ (i.e. there exists a value v_i which is a function of v_j) with $(v_i(v_j), v_j) \in R_{ij}$ such that $\forall x_k \in X \setminus \{x_i, x_j\}, \forall v_k \in \mathcal{D}(x_k)$ with $(v_j, v_k) \in R_{jk}$, we have $(v_i(v_j), v_k) \in R_{ik}$. Define s as follows: $s(x_m) = s'(x_m)$ ($x_m \in X \setminus \{x_i\}$) and $s(x_i) = v_i(s'(x_j))$. The assignment $\langle x_i, v_i(s'(x_j)) \rangle$ is compatible with $\langle x_j, s'(x_j) \rangle$ (by definition of $v_i(s'(x_j))$) and is compatible with all of the assignments $\langle x_m, s'(x_m) \rangle$ ($x_m \in X \setminus \{x_i\}$) again by definition of $v_i(s'(x_j))$ since $(s'(x_j), s'(x_m)) \in R_{jm}$. Hence s is a solution to I .

It is easily verified that this proof is valid even in the special case $X = \{x_i, x_j\}$. ■

The proof of the following proposition can be found in Appendix C where we give an optimised algorithm, making use of appropriate data structures, to apply this variable-elimination rule until convergence.

Proposition 3 *Variable eliminations by the triangle property can be applied until convergence in $O(\text{end}^3)$ time and $O(\text{end}^2)$ space.*

5. From broken triangles to broken polyhedra

The broken-triangle property is a property of the microstructure of instances of the binary CSP (Constraint Satisfaction Problem) which when satisfied allows either value merging (Cooper, Duchain, El Mouelhi, Escamocher, Terrioux, & Zanuttini, 2016a), variable elimination or the definition of a tractable class (Cooper et al., 2010). In this section, we generalise the notion of broken triangle to broken polyhedron, which allows us to define rules for variable elimination parameterised by the dimension k of the polyhedron.

We begin by recalling the definition of the broken-triangle property (BTP) (Cooper et al., 2010).

Definition 6 *Let $I = \langle X, \mathcal{D}, R \rangle$ be a binary CSP instance. A pair of values $v'_k, v''_k \in \mathcal{D}(x_k)$ satisfies BTP if for each pair of variables (x_i, x_j) (with $i, j \neq k$), $\forall v_i \in \mathcal{D}(x_i), \forall v_j \in \mathcal{D}(x_j)$, if*

- $(v_i, v_j) \in R_{ij}$,
- $(v_i, v'_k) \in R_{ik}$ and
- $(v_j, v''_k) \in R_{jk}$,

then

- $(v_i, v''_k) \in R_{ik}$ or
- $(v_j, v'_k) \in R_{jk}$.

A variable x_k satisfies BTP if each pair of values of $\mathcal{D}(x_k)$ satisfies BTP. If I is equipped with an order $<$ on its variables, then I satisfies BTP for the variable order $<$ if each variable x_k satisfies BTP in the sub-instance of I restricted to the variables x_i such that $x_i \leq x_k$.

If $(v_i, v_j) \in R_{ij}, (v_i, v'_k) \in R_{ik}, (v_j, v''_k) \in R_{jk}, (v_i, v''_k) \notin R_{ik}$ and $(v_j, v'_k) \notin R_{jk}$ (as shown in Figure 4 with solid/broken lines joining compatible/incompatible values), then the quadruple (v'_k, v_i, v_j, v''_k) constitutes a *broken triangle* on x_k . I satisfies the BTP on x_k if no broken triangles occur on x_k .

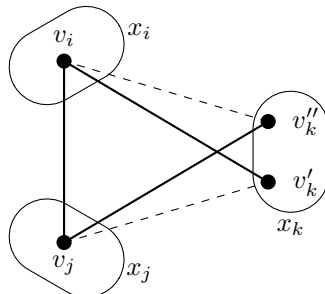


Figure 4: A broken triangle (v'_k, v_i, v_j, v''_k) .

Any pair of values v'_k, v''_k that satisfy BTP can be merged without changing the satisfiability of the instance (Cooper et al., 2016a), where the new merged value is compatible with all values compatible with at least one of the two old values v'_k, v''_k . Furthermore, in an arc-consistent instance any variable that satisfies BTP can be eliminated without changing the satisfiability of the instance (Cooper et al., 2010). As a direct consequence of this, if an arc-consistent instance I satisfies BTP for some variable ordering, then I can be solved in polynomial time by successive elimination of all variables: moreover viewed as a decision problem, I can, in fact, be solved by arc consistency (Cooper et al., 2010) without knowledge of the variable ordering for which BTP holds.

Examples of the BTP var-elim rule include a variable x_m which is only constrained by one other variable in an arc-consistent instance or a variable x_m with a domain of size at most two in a path-consistent instance (Cooper et al., 2010). In this paper we consider generalisations of BTP which allow the definition of stronger variable-elimination rules.

The presence of some broken triangles on a given variable does not preclude value-merging or variable-elimination (while leaving satisfiability invariant): new solutions will not be introduced if the broken triangles lack support on some set of other variables (Cooper, El Mouelhi, & Terrioux, 2016b; El Mouelhi, 2018; Naanaa, 2016). Unfortunately, the search for lack-of-support variables for each broken triangle may render such techniques prohibitively expensive in terms of time complexity (since, in the worst case, the number of broken triangles is $\Theta(n^3 d^4)$). Other generalisations of BTP require levels of consistency, such as strong path consistency, which may change positive edges into negative edges; this has the disadvantage of possibly introducing new broken triangles besides the extra memory required to store new binary constraints (Cooper, Jégou, & Terrioux, 2015; Naanaa, 2013).

We now generalise the notion of broken triangle to broken polyhedron. In Section 6 we show that this notion can be used to define variable elimination rules that are stronger than BTP. A broken triangle is a broken polyhedron of dimension 2. We now define a broken k -dimensional polyhedron for $k \geq 2$.

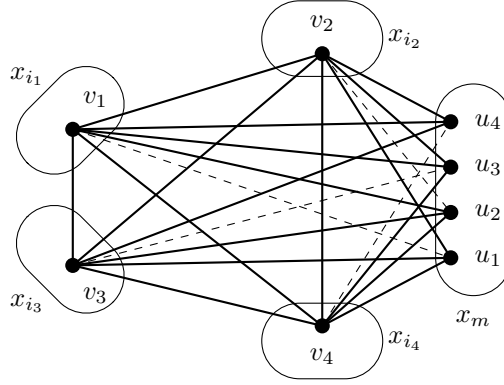


Figure 5: A broken 4-dimensional polyhedron on variable x_m .

Definition 7 A broken k -dimensional polyhedron on x_m consists of

- a consistent assignment $\langle v_1, \dots, v_k \rangle$ to distinct variables $\langle x_{i_1}, \dots, x_{i_k} \rangle$ (where each x_{i_j} ($j = 1, \dots, k$) is distinct from x_m),
- k distinct values $u_1, \dots, u_k \in D_m$,

such that

- $\forall j \in \{1, \dots, k\}, (v_j, u_j) \notin R_{i_j m}$,
- $\forall h, j \in \{1, \dots, k\},$ if $h \neq j$ then $(v_h, u_j) \in R_{i_j m}$,

The assignment $\langle v_1, \dots, v_k \rangle$ to variables $\langle x_{i_1}, \dots, x_{i_k} \rangle$ is known as the base of the broken polyhedron, and each assignment $\langle x_m, u_j \rangle$ ($j = 1, \dots, k$) is an apex. Any edge between a base point and an apex is a side of the broken tetrahedron.

A broken 4-dimensional polyhedron is shown in Figure 5. A broken triangle (Figure 4) is a broken 2-dimensional polyhedra. In the following sections we show that broken k -dimensional polyhedra allow us to define novel variable-elimination rules and tractable classes.

6. First-order rules for variable-elimination by broken polyhedra

The broken-triangle property (BTP) has been generalised to the $\forall\exists$ BTP rule for variable elimination which allows us to eliminate more variables (Cooper, 2014) than BTP. Eliminating a variable satisfying the $\forall\exists$ BTP rule is strictly stronger than the BTP rule. This is demonstrated by the fact that $\forall\exists$ BTP, but not BTP, subsumes the rule that allows us to eliminate a variable x_m when an assignment to x_m is compatible with all assignments to all other variables. Another generic example is when all occurrences of the broken-triangle pattern on variable x_m occur on pairs of values $v_m, v'_m \in S \subset \mathcal{D}(x_m)$ and each assignment v_i to each other variable $x_i \neq x_m$ has a support at x_m in $\mathcal{D}(x_m) \setminus S$.

We first given the definition of the $\forall\exists$ broken-triangle property, in order to generalise it to k dimensions.

Definition 8 A binary CSP instance satisfies the $\forall\exists$ broken-triangle property on variable x_m if for all $i_1 \neq m$, for all $v_1 \in \mathcal{D}(x_{i_1})$, there exists $v_m \in \mathcal{D}(x_m)$ such that

1. $\langle v_1, v_m \rangle$ is a consistent assignment to variables $\langle x_{i_1}, x_m \rangle$, and
2. for all $i_2 \notin \{i_1, m\}$, for all $v_2 \in \mathcal{D}(x_{i_2})$, there is no broken triangle on x_m with base the assignment $\langle v_1, v_2 \rangle$ to variables $\langle x_{i_1}, x_{i_2} \rangle$ and with an apex $\langle x_m, v_m \rangle$.

We now generalise the $\forall\exists$ BTP rule for variable elimination to the case of broken polyhedra of any dimension $k \geq 2$. When $k = 2$ the following definition coincides with Definition 8 of the $\forall\exists$ BTP rule.

Definition 9 A binary CSP instance satisfies the $\forall\exists$ broken k -dimensional polyhedron property on variable x_m if for all distinct $i_1, \dots, i_{k-1} \neq m$, for all consistent assignments $\langle v_1, \dots, v_{k-1} \rangle$ to variables $\langle x_{i_1}, \dots, x_{i_{k-1}} \rangle$, there exists $v_m \in \mathcal{D}(x_m)$ such that

1. $\langle v_1, \dots, v_{k-1}, v_m \rangle$ is a consistent assignment to variables $\langle x_{i_1}, \dots, x_{i_{k-1}}, x_m \rangle$, and
2. for all $i_k \notin \{i_1, \dots, i_{k-1}, m\}$, for all $v_k \in \mathcal{D}(x_{i_k})$, there is no broken k -dimensional polyhedron on x_m with base the assignment $\langle v_1, \dots, v_k \rangle$ to variables $\langle x_{i_1}, \dots, x_{i_k} \rangle$ and with an apex $\langle x_m, v_m \rangle$.

The first condition of Definition 9 is a k -consistency condition on variable x_m with respect to all other variables (Lecoutre, 2009). The second condition guarantees (as we will show below) that this k -consistency condition is sufficient for any consistent assignment to the variables $X \setminus \{x_m\}$ to be extendible to a consistent assignment to all variables.

Notation: Given a binary CSP instance I on variables $X = \{x_1, \dots, x_n\}$, we denote by I_{-m} the sub-instance of I on variables $X \setminus \{x_m\}$. Similarly, I_{-jm} denotes the sub-instance of I on variables $X \setminus \{x_j, x_m\}$.

We can observe that if I satisfies the $\forall\exists$ broken k -dimensional polyhedron property on x_m then I_{-j} also satisfies the $\forall\exists$ broken k -dimensional polyhedron property on x_m for any $j \neq m$.

Theorem 3 The $\forall\exists$ broken k -dimensional polyhedron property is a sol-var-elim condition in binary CSP instances I with at least k variables.

Proof: Let I be a binary CSP instance I on $n \geq k$ variables which satisfies the $\forall\exists$ broken k -dimensional polyhedron property on variable x_m . It is sufficient to show that any solution s to I_{-m} can be extended to a solution to I . We will show this by induction on n . That a solution for I can be generated in polynomial time will follow immediately since s does not need to be modified, just extended by one of the at most d possible values for x_m .

If $n = k$, then the fact that any solution to I_{-m} can be extended to a solution to I follows directly from the definition of the $\forall\exists$ broken k -dimensional polyhedron property. So, to complete the proof by induction, we suppose that any solution to I_{-m} can be extended to a solution to I holds for instances I with $n = p \geq k$ variables and we will show that this also holds for instances with $p + 1$ variables.

For notational convenience and without loss of generality, we can assume that $m = p + 1$. Let $s := \langle v_1, \dots, v_p \rangle$ be a solution to I_{-m} . To complete the proof, it suffices to show that s can be extended to a solution to I . For each $j = 1, \dots, p$, consider the p -variable instance I_{-j} . Clearly, $s_j :=$

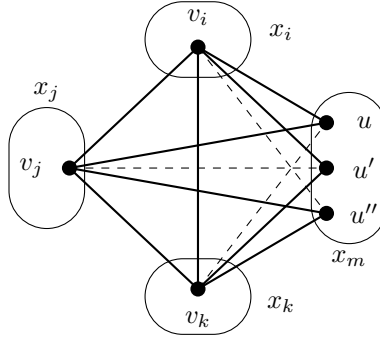


Figure 6: A broken tetrahedron.

$\langle v_1, \dots, v_{j-1}, v_{j+1}, \dots, v_p \rangle$ is a solution to I_{-j_m} . Since I_{-j} has p variables, by our inductive hypothesis, each s_j ($j = 1, \dots, p$) can be extended to a solution $t_j = \langle v_1, \dots, v_{j-1}, v_{j+1}, \dots, v_p, u_j \rangle$ to I_{-j} . Note that t_j assigns u_j to x_m . Suppose that for some $j \in \{1, \dots, p\}$, $(v_j, u_j) \in R_{j_m}$. Then $\langle v_1, \dots, v_p, u_j \rangle$ is a solution to I and we are done. So, we only need consider the case in which $\forall j \in \{1, \dots, p\}$, $(v_j, u_j) \notin R_{j_m}$. Note that the values u_j ($j = 1, \dots, p$) must all be distinct.

Consider $\langle v_1, \dots, v_{k-1} \rangle$; this is a consistent assignment to variables $\langle x_1, \dots, x_{k-1} \rangle$. Thus, since I satisfies the $\forall \exists$ broken k -dimensional polyhedron property on variable x_m , there exists $v_m \in \mathcal{D}(x_m)$ such that (1) $\langle v_1, \dots, v_{k-1}, v_m \rangle$ is a consistent assignment to the variables $\langle x_1, \dots, x_{k-1}, x_m \rangle$ and (2) for all $h \in \{k, \dots, p\}$, there is no broken k -dimensional polyhedron on x_m with base the assignment $\langle v_1, \dots, v_{k-1}, v_h \rangle$ to variables $\langle x_1, \dots, x_{k-1}, x_h \rangle$ and with an apex $\langle x_m, v_m \rangle$. Note that v_h is the value assigned by s to x_h . Observe that $\langle v_1, \dots, v_{k-1}, v_h \rangle$ is a consistent assignment to variables $\langle x_1, \dots, x_{k-1}, x_h \rangle$ and that u_1, \dots, u_{k-1}, v_m are distinct values (since $\forall j \in \{1, \dots, p\}$, $(v_j, u_j) \notin R_{j_m}$ but $(v_j, v_m) \in R_{j_m}$, and we have already seen above that the u_j ($j = 1, \dots, p$) are distinct). Furthermore, by the definition of the u_j , we have $(v_h, u_j) \in R_{hm}$ for $j = 1, \dots, k-1$, and $(v_\ell, u_j) \in R_{\ell m}$ if and only if $\ell \neq j$ for $\ell, j \in \{1, \dots, k-1\}$. By the definition of v_m (condition (1) above), we have $(v_j, v_m) \in R_{j_m}$ for $j = 1, \dots, k-1$. Since there is no broken k -dimensional polyhedron on x_m with base the assignment $\langle v_1, \dots, v_{k-1}, v_h \rangle$ to variables $\langle x_1, \dots, x_{k-1}, x_h \rangle$ and with apex $\langle x_m, v_m \rangle$, we must have $(v_h, v_m) \in R_{hm}$. Since this is true for all $h \in \{k, \dots, p\}$, it follows that $\langle v_1, \dots, v_p, v_m \rangle$ is a solution to I . This completes the proof by induction. ■

7. Faster variable-elimination based on broken polyhedra

The $\forall \exists$ broken k -dimensional polyhedron property is interesting from a theoretical point of view. However, from a practical point of view, the time complexity of detecting whether variables can be eliminated is likely to be prohibitive. Indeed, for $k = 3$, a naive exhaustive search for broken tetrahedra (i.e. broken 3-dimensional polyhedra) has time complexity $\Theta(n^4 d^6)$. Various intermediate rules of varying strength and time complexity exist between $\forall \exists$ BTP and the $\forall \exists$ broken tetrahedron property; we choose to concentrate on rules which can be tested in the same worst-case time complexity as $\forall \exists$ BTP but which are strictly stronger.

First, observe that a broken tetrahedron, as shown in Figure 6 with base the assignment $\langle v_i, v_j, v_k \rangle$ to variables $\langle x_i, x_j, x_k \rangle$ and apexes $u, u', u'' \in \mathcal{D}(x_m)$, contains three broken triangles on x_m : (u', v_i, v_j, u'') , (u, v_i, v_k, u'') and (u, v_j, v_k, u') . Thus, the incompatible pairs (shown as broken

lines in Figure 6) (v_i, u'') and (v_j, u') each occur in (at least) two broken triangles, while the compatible pairs (shown as solid lines in the figure) (v_i, v_j) , (v_i, u) and (v_j, u) each occur in (at least) one broken triangle.

Definition 10 Let I be a binary CSP instance with $v_i \in \mathcal{D}(x_i)$ and $u \in \mathcal{D}(x_m)$. The broken-triangle degree (BT degree) of the pair of assignments $\langle v_i, u \rangle$ to $\langle x_i, x_m \rangle$ is the number of distinct variables x_j ($j \neq i, m$) such that $\exists v_j \in \mathcal{D}(x_j), \exists u' \in \mathcal{D}(x_m)$ such that there is a broken triangle with base the assignment $\langle v_i, v_j \rangle$ to $\langle x_i, x_j \rangle$ and apexes $u, u' \in \mathcal{D}(x_m)$.

For example, if I is exactly the instance shown in Figure 6, then the BT degree of (v_i, u'') and of (v_j, u') is two, and the BT degree of (v_i, u) and of (v_j, u) is one. Although (v_i, v_j) is the base of a broken triangle, its BT degree is zero according to Definition 10, since neither v_i nor v_j is the apex of a broken triangle. Note that any edge, whether positive or negative, which links the base to an apex of a broken triangle has BT degree at least one.

Definition 11 A consistent assignment $\langle v_i, v_j \rangle$ to variables $\langle x_i, x_j \rangle$ is 3-safe on variable x_m if for all broken triangles (u', v_i, v_j, u'') , the BT degree of (v_i, u'') is one or the BT degree of (v_j, u') is one.

Note that $\langle v_i, v_j \rangle$ is trivially 3-safe on x_m if there is no broken triangle on x_m whose base is $\langle v_i, v_j \rangle$. We can now define a new variable-elimination property.

Definition 12 A binary CSP instance satisfies the BT-degree property on variable x_m if for all distinct $i, j \neq m$, for all $v_i \in \mathcal{D}(x_i), v_j \in \mathcal{D}(x_j)$ such that $\langle v_i, v_j \rangle$ is a consistent assignment to variables $\langle x_i, x_j \rangle$, there exists $v_m \in \mathcal{D}(x_m)$ such that

1. $\langle v_i, v_j, v_m \rangle$ is a consistent assignment to variables $\langle x_i, x_j, x_m \rangle$, and
2. either (v_i, v_j) is 3-safe on x_m or (v_i, v_m) has BT degree zero or (v_j, v_m) has BT degree zero.

The first condition in Definition 12 guarantees path consistency on variable x_m with respect to all other pairs of variables (Lecoutre, 2009), whereas the second condition guarantees the absence of a broken 3-dimensional polyhedron with base points v_i, v_j and apex v_m .

Theorem 4 The BT-degree property is a sol-var-elim condition in binary CSP instances I with at least 3 variables.

Proof: It suffices, by Theorem 3, to show that a binary CSP instance I that satisfies the BT-degree property on x_m necessarily satisfies the $\forall\exists$ broken 3-dimensional polyhedron property on x_m . But this is immediate by the discussion above since

1. pairs of assignments which are 3-safe on x_m cannot be part of the base of a broken tetrahedron on x_m , by the remark before Definition 10,
2. pairs of assignments which have BT degree zero cannot be the side (i.e. edge between a base point and an apex) of a broken tetrahedron.

■

Variable eliminations by the BT-degree property may propagate. In the instance I of Figure 7(a) (where in this figure pairs of values not joined by a line are assumed to be incompatible), the variables x_1, x_2, x_3 cannot be eliminated by Theorem 4 since the consistent assignments (v_3, v'_4) , (v'_1, v_4) and (v'_1, v''_4) have no support, respectively, at x_1, x_2 and x_3 . However, we can eliminate x_4 by Theorem 4, since v_4 is a support at x_4 of any consistent assignment (v_i, v_j) to any pair of variables (x_i, x_j) ($1 \leq i < j \leq 3$) and in each case (v_i, v_4) has BT degree 0 (since v_4 is consistent with all assignments to all other variables). Eliminating x_4 then produces an instance I' (shown in Figure 7(b)) in which x_1 can now be eliminated by Theorem 4 since the only consistent assignment (v_2, v_3) to (x_2, x_3) can be extended to the consistent assignment (v_2, v_3, v_1) to (x_2, x_3, x_1) and (v_2, v_1) has BT degree 0.

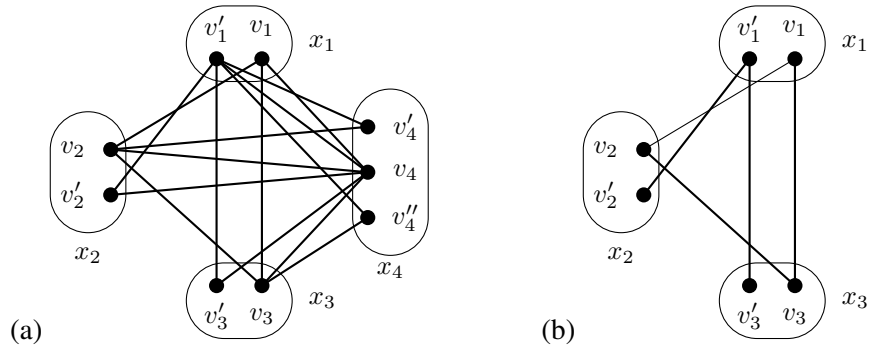


Figure 7: (a) x_1 cannot be eliminated (by the BT-degree property) in this instance I , (b) x_1 can be eliminated in this instance I' which results after elimination of x_4 from I .

The proof of the following proposition can be found in Appendix D where we give an optimised algorithm, using appropriate data structures, to apply this variable-elimination rule until convergence.

Proposition 4 *The BT-degree variable-elimination property can be applied until convergence in $O(end^3)$ time and $O(end^2)$ space.*

We can compare this with the $\forall\exists$ BTP property (Definition 9 with $k = 2$) which can also be applied until convergence in $O(end^3)$ time and $O(end^2)$ space, as shown in Appendix E. Eliminating variables by the broken triangle property, which is subsumed by the $\forall\exists$ BTP property, also has time complexity $O(end^3)$ (Cooper et al., 2010). Thus the BT-degree property is comparable with the $\forall\exists$ BTP property in terms of computational complexity. The following proposition shows that it is at least as powerful in terms of number of variables eliminated.

Proposition 5 *If a binary CSP instance satisfies $\forall\exists$ BTP on variable x_m then it satisfies the BT-degree property on variable x_m .*

Proof: The $\forall\exists$ BTP property says that for all $i \neq m$, for all $v_i \in \mathcal{D}(x_i)$, $\exists d(v_i) \in \mathcal{D}(x_m)$ (i.e. there exists a value d which is a function of v_i) such that $(v_i, d(v_i)) \in R_{im}$ and the BT degree of $(v_i, d(v_i))$ is zero. To show that this implies the BT-degree property on variable x_m , consider any pair of distinct variables x_i, x_j ($i, j \neq m$) and any values $v_i \in \mathcal{D}(x_i)$, $v_j \in \mathcal{D}(x_j)$ such that $(v_i, v_j) \in R_{ij}$. We need to show that there exists $v_m \in \mathcal{D}(x_m)$ such that

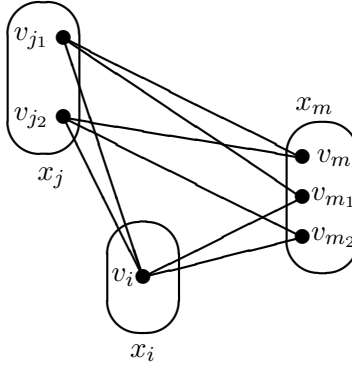


Figure 8: An instances in which x_m can be eliminated by the BT-degree property but not by the $\forall\exists$ BTP property.

1. $(v_i, v_m) \in R_{im}$ and $(v_j, v_m) \in R_{jm}$, and
2. (v_i, v_j) is 3-safe on x_m or the BT degree of (v_i, v_m) is zero or the BT degree of (v_j, v_m) is zero.

If $(v_j, d(v_i)) \in R_{jm}$, then $v_m = d(v_i)$ satisfies the first condition and the BT degree of (v_i, v_m) is zero. If $(v_i, d(v_j)) \in R_{im}$, then $v_m = d(v_j)$ satisfies the first condition and the BT degree of (v_j, v_m) is zero. It is not possible to have both $(v_j, d(v_i)) \notin R_{jm}$ and $(v_i, d(v_j)) \notin R_{im}$ since, in this case, $(d(v_i), v_i, v_j, d(v_j))$ would be a broken triangle, contradicting the fact that the BT degree of $(v_i, d(v_i))$ is zero. We can therefore conclude that in all cases, the BT-degree property is satisfied on variable x_m . ■

In fact, the BT-degree property strictly subsumes the $\forall\exists$ BTP property, as illustrated by the instance in Figure 8. In this instance, x_m cannot be eliminated by $\forall\exists$ BTP since for $v_i \in \mathcal{D}(x_i)$, there is no $v \in \mathcal{D}(x_m)$ with $(v_i, v) \in R_{im}$ and such that the BT degree of (v_i, v) is zero (because of the broken triangles $(v_{m1}, v_i, v_{j2}, v_m)$ and $(v_{m2}, v_i, v_{j1}, v_m)$). On the other hand, x_m can be eliminated by the BT-degree property since for both of the assignments (v_i, v_{j_k}) ($k = 1, 2$) to (x_i, x_j) , there is a value $v_{m_k} \in \mathcal{D}(x_m)$ with $(v_i, v_{m_k}) \in R_{im}$, $(v_{j_k}, v_{m_k}) \in R_{jm}$ and such that the BT degree of (v_{j_k}, v_{m_k}) is zero.

8. Theoretical comparison between different variable-elimination rules

The notion of rank introduced by Naanaa (Naanaa, 2013) is closely related to the absence of broken k -dimensional polyhedra. Indeed, in a binary CSP instance, a variable x_m has rank $k - 1$ if there is no broken k -dimensional polyhedron on x_m . Naanaa showed that a variable x_m with rank at most $k - 1$ in a binary CSP instance which is directional strong k -consistent (Lecoutre, 2009) (according to an order which places x_m last) can be eliminated while leaving the satisfiability of the instance invariant. This is subsumed by the $\forall\exists$ broken k -dimensional polyhedron property (Definition 9) since the latter does not require the absence of all broken k -dimensional polyhedra (Cooper & Zivny, 2017).

Following an orthogonal approach, it has recently been shown that singleton arc consistency (Lecoutre, 2009) solves instances that do not contain a pattern (known as $Q1$) made up of a subset of the edges of a broken tetrahedron (Carbonnel, Cohen, Cooper, & Zivny, 2018).

Another family of variable-elimination rules is m -fBTP (El Mouelhi, 2018), for $m \geq 1$, which extends BTP by allowing broken triangles that do not have a support at some subset of variables of size m .

Definition 13 *Let $I = \langle X, \mathcal{D}, R \rangle$ be a binary CSP instance. A pair of values $v'_k, v''_k \in \mathcal{D}(x_k)$ satisfies 1-fBTP if for each broken triangle (v'_k, v_i, v_j, v''_k) with $v_i \in \mathcal{D}(x_i)$, $v_j \in \mathcal{D}(x_j)$, there is at least one variable $x_\ell \in X \setminus \{x_i, x_j, x_k\}$ such that $\forall v_\ell \in \mathcal{D}(x_\ell)$, if $(v_i, v_\ell) \in R_{i\ell}$ then $(v_j, v_\ell) \notin R_{j\ell}$. In this case, we say that x_ℓ is a support variable for the broken triangle (v'_k, v_i, v_j, v''_k) . A variable $x_k \in X$ satisfies 1-fBTP if each pair of values $v'_k, v''_k \in \mathcal{D}(x_k)$ satisfies 1-fBTP.*

The variable x_m of the instance shown in Figure 6 does not satisfy 1-fBTP because there is no support variable for the broken triangle (u, v_i, v_k, u'') . The concept of support variable can be extended to a set of m variables to obtain the definition of m -fBTP (El Mouelhi, 2018).

In the rest of the paper, we use the notation $(\exists, \text{DE-})$ snake in the following sense: a statement is true for $(\exists, \text{DE-})$ snake property if it is true for both the \exists snake property and the DE-snake property. We now compare theoretically the four variable-elimination rules: the $(\exists, \text{DE-})$ snake property, the triangle property, $\forall\exists$ BTP and 1-fBTP. Two variable-elimination rules are incomparable if neither is subsumed by the other.

Proposition 6 *The following four variable-elimination rules are all pairwise incomparable: the $(\exists, \text{DE-})$ snake property, the triangle property, $\forall\exists$ BTP and 1-fBTP.*

Proof: Figure 9 shows four binary CSP instances. In this figure, compatible values are joined by lines, and hence incompatibility is represented by the absence of a line. In the instance shown in Figure 9(a), variable x_i can be eliminated by the \exists snake property, the triangle property or $\forall\exists$ BTP, but not by 1-fBTP. Indeed, a variable x_i for which there exists $v''_i \in \mathcal{D}(x_i)$ compatible with all values for all other variables (as is the case for x_i in Figure 9(a)) can always be eliminated by the \exists snake (and also the DE-snake) property, the triangle property or $\forall\exists$ BTP, but not necessarily by 1-fBTP. In this example, there is a broken triangle (v'_i, v_k, v_j, v_i) , shown in red, and (trivially) no other variable on which this broken triangle does not have a support, so x_i cannot be eliminated by 1-fBTP.

In the instance shown in Figure 9(b), there is no broken triangle on variable x_i , so it can be eliminated by $\forall\exists$ BTP or 1-fBTP, but not by the $(\exists, \text{DE-})$ snake property nor the triangle property. It is easily verified that the snake pattern (Figure 1) occurs on each value $u \in \mathcal{D}(x_i)$ for some other variable x_m and for each positive edge uv with $v \in \mathcal{D}(x_m)$ (the snake patterns are represented by three different colours with the negative edges of the pattern shown as dashed lines).

In the instance shown in Figure 9(c), variable x_i can be eliminated by the \exists snake property (and hence also the DE-snake property which subsumes the \exists snake property), but none of the triangle property, $\forall\exists$ BTP or 1-fBTP. The \exists snake property is satisfied since the snake pattern (Figure 1) does not occur on $v_i \in \mathcal{D}(x_i)$. The triangle property is not satisfied on x_i since the open-triangle pattern shown in Figure 3 occurs on $v_j \in \mathcal{D}(x_j)$ and on $v_k \in \mathcal{D}(x_k)$. The broken triangles (v_i, v_j, v_k, v'_i) and (v_i, v'_k, v'_j, v'_i) , respectively shown in blue and red, prevent elimination of x_i by $\forall\exists$ BTP or 1-fBTP.

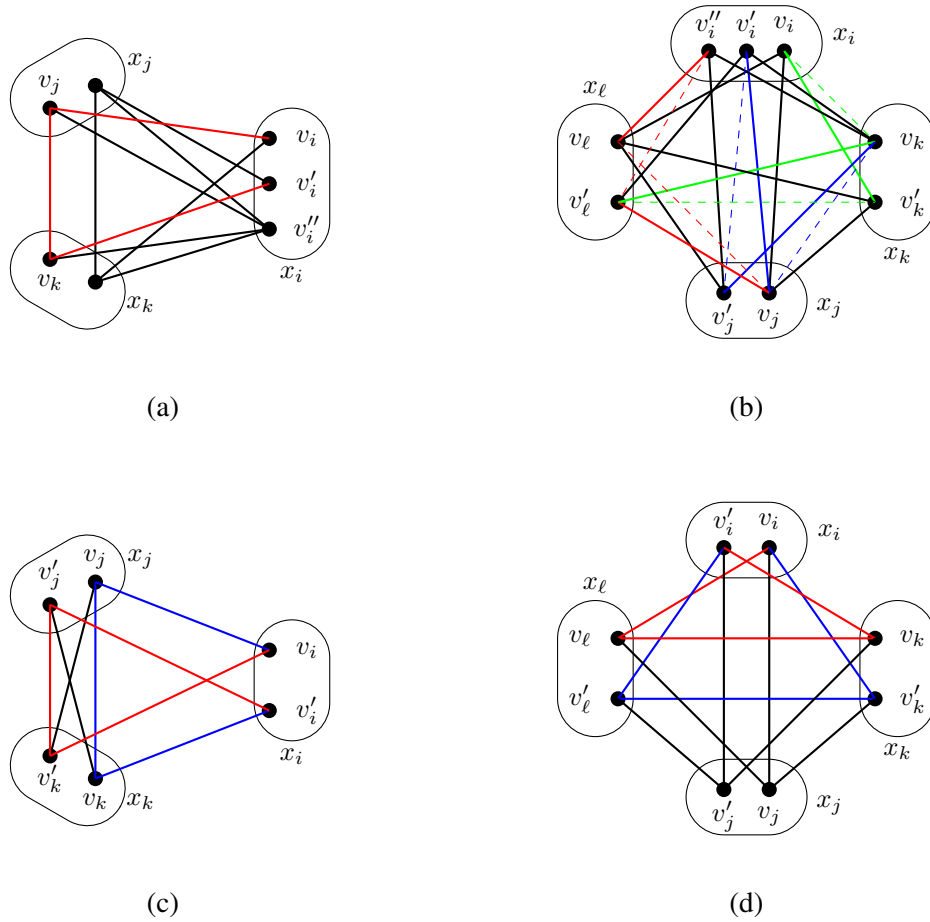


Figure 9: Examples of instances to demonstrate the incomparability of the different variable-elimination rules.

In the instance shown in Figure 9(d), variable x_i can be eliminated by the triangle property or 1-fBTP, but not by $\forall\exists$ BTP nor the $(\exists, \text{DE-})$ snake property. The triangle property is satisfied on x_i since, for all $v \in \mathcal{D}(x_j)$ no open-triangle pattern (illustrated in Figure 3) occurs on v . The broken triangles (v_i, v_ℓ, v_k, v'_i) and $(v'_i, v'_\ell, v'_k, v_i)$, respectively shown in red and blue, prevent elimination of x_i by $\forall\exists$ BTP; but these broken triangles have a support variable x_j which allows x_i to be eliminated by 1-fBTP.

It can easily be verified that all combinations are covered by these four examples: for any two distinct rules, rule1 and rule2, among the $(\exists, \text{DE-})$ snake property, the triangle property, $\forall\exists$ BTP and 1-fBTP, there is an example instance in Figure 9 such that rule1 eliminates x_i but rule2 does not. ■

Proposition 5 tells us that the BT-degree property subsumes $\forall\exists$ BTP. On the other hand, as we now show, it is incomparable with the three other properties.

Proposition 7 *The BT-degree property is incomparable with each of the following variable elimination properties: the $(\exists, \text{DE-})$ snake property, the triangle property and 1-fBTP.*

Proof: The proof is identical to the proof of Proposition 6, since the instances shown in Figure 9 in which variable x_i satisfies the BT-degree property are exactly the same instances in which x_i satisfies $\forall\exists\text{BTP}$. In particular, x_i does not satisfy the BT-degree property in Figure 9(c) (respectively Figure 9(d)) since (v_j, v_k) (respectively (v_ℓ, v_k)) cannot be extended to a consistent assignment for x_i . Proposition 5 tells us that x_i can be eliminated by BT-degree property in the instances in Figure 9(a) and Figure 9(b) since it can be eliminated by the weaker property $\forall\exists\text{BTP}$. ■

Another important generic example is the case in which x_i is constrained by a single other variable x_j . Such a variable x_i can always be eliminated (remembering that eliminating a variable means first deleting from $\mathcal{D}(x_j)$ all values v_j with no support at x_i). In this case, variable x_i can be eliminated by any of the triangle property, $\forall\exists\text{BTP}$ or 1-fBTP, but not necessarily by the $(\exists, \text{DE-})$ snake property.

9. Experimental results

In this section, we study the practical interest of some variable elimination rules, namely the BT-degree property, the \exists snake rule, the DE-snake rule and the triangle property. For each rule, we assess its ability to eliminate variables and its impact on solving efficiency. First, we describe the experimental protocol we used.

9.1 Experimental protocol

We considered all the binary instances from the 2008 International CP Competition¹ and we discarded those whose inconsistency is detected by enforcing arc-consistency. By so doing, we obtained a benchmark of 3,557 CSP instances. These instances have between 3 and 5,000 variables whose initial domains have between 2 and 10,000 values. The number of constraints varies from 3 to 124,750. Constraints are defined in extension or in intension. For example, among these instances, we can find frequency allocation problems or graph colouring instances.

Regarding the variable elimination algorithms, for each rule, we first enforced arc-consistency and then eliminated those variables having a singleton domain. Note that, in an arc-consistent instance, singleton-domain variables would be eliminated by all four of the rules we are comparing. Then we applied the considered elimination rule until convergence (i.e. a fixpoint is reached at which no more eliminations are possible by this rule). To do this, we consider a set of variables containing all the candidates for elimination. Initially, this set contains all the variables having a non-singleton domain. For each candidate variable, we check whether the rule applies. If so, the variable is eliminated and all its neighbours are added to the candidate set. Two variables are neighbours if they constraint each other. Checking whether the rule applies is performed as described in Section 7 for the BT-degree property, or by a naive approach for the \exists snake rule, the DE-snake rule and the triangle property. We used more naive algorithms than those which are optimised for worst-case time complexity. The algorithms in the Appendix (which use data structures to reduce worst-case time complexity) are given for their theoretical rather than practical interest, since our first concern in these experimental trials was to estimate and compare the number of variable eliminations that can be achieved by each rule. In order to solve CSP instances, we used the state-of-the-art algorithm MAC+RST+NG (Lecoutre, Sais, Tabary, & Vidal, 2007). We

1. <http://www.cril.univ-artois.fr/CPAI08>

Table 1: Number of instances for which the elimination process finishes, runs out of time or memory or is able to eliminate at least one variable.

	BT-degree	\exists snake	DE-snake	triangle
# Processed	2,056	3,449	3,371	3,420
# Timeout	513	108	184	137
# Memory-out	988	0	0	0
# Elim. instances	507	786	836	1,313

exploit a geometric restart policy based on the number of allowed backtracks. Initially, the number of allowed backtracks is set to 100 and the increasing factor to 1.1. The search was guided by the dom/wdeg variable heuristic (Boussemart, Hemery, Lecoutre, & Sais, 2004). All the algorithms are written in C++ in our own library.

The experiments were performed on Dell PowerEdge M620 blade servers with Intel Xeon E5-2609 2.4 GHz processors. We allotted 30 minutes and 16 GB of memory for each elimination process while, for the solving process, the timeout was set to one hour.

9.2 Ability to eliminate variables

In this part, we assess the practical behaviour of the considered elimination rules and their ability to eliminate variables. Table 1 provides the number of instances for which the elimination process finishes, runs out of time or memory or is able to eliminate at least one variable. Clearly, the $(\exists, \text{DE-})$ snake rules and the triangle property are able to process more instances than the BT-degree property. Such a result was foreseeable since the algorithm we used for the latter property has worse time and space complexities than the two others. These complexities also explain why the elimination process based on the BT-degree property runs out of time or memory. However, despite this, the BT-degree property still succeeds in processing about 58% of the instances. Moreover, as we can see in Figure 10, which depicts the cumulative number of processed instances with respect to elapsed time, it processes about 90% of the treated instances in less than one minute. This ratio is close to the one obtained for the triangle property, while the \exists snake rule (respectively the DE-snake rule) can be applied for more than 95% (resp. 92%) of the treated instances in the same time. An obvious conclusion we can draw from Figure 10 is that a much shorter time-out would not have greatly reduced the number of variable eliminations by any of the rules. Indeed, the fact that the curves all flatten out fairly quickly indicates that more efficient algorithms would not have detected a significantly larger number of variable eliminations within any given time-out period.

Now, if we investigate the ability to eliminate variables, the triangle property is the most interesting rule in the sense that it is able to eliminate variables in more instances. The triangle rule eliminates at least one variable in 1,313 instances. Figure 11 provides a comparison of the percentage of eliminated variables for each instance and for each pair of elimination rules. First, we can remark that these comparisons are consistent with the theoretical results we provided previously. In particular, they illustrate the fact that the DE-snake rule subsumes the \exists snake rule and any other pair of elimination rules are incomparable. For example, if we compare the $(\exists, \text{DE-})$ snake rules and the BT-degree property, we clearly see that they are incomparable since there exist instances for which some variables are eliminated by the first elimination rules and not by the second and conversely. Moreover, we can note that the number of the instances for which the $(\exists, \text{DE-})$ snake rules

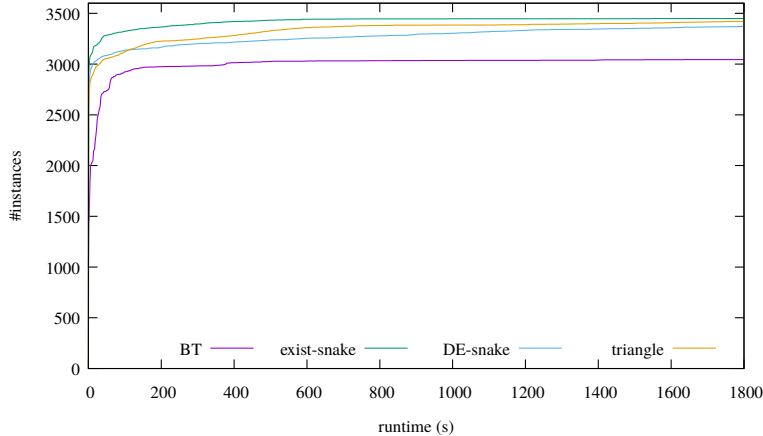


Figure 10: Cumulative number of processed instances with respect to elapsed time.

eliminate more variables is close to the respective number for the BT-degree property. In contrast, if we compare the triangle property with the BT-degree property, the triangle property turns out to be more effective for eliminating variables, even if the two rules are incomparable. Indeed, there exist only a few instances for which the BT-degree property eliminates more variables than the triangle property. Then, if we compare the triangle property with the \exists snake rule, we can draw the same conclusion. Finally, the comparison between the triangle property and the DE-snake rule seems to be more less clear. However, we can note that for a significant number of instances, the triangle property is able to eliminate variables while the DE-snake rule eliminates none. So, it turns out that the triangle property appears to be the best elimination rule with respect to the ability to eliminate variables.

Further experiments would be needed to identify which type of variables are more likely to be eliminated by each rule. If it turns out that such variables have some easily identifiable characteristic, such as a small number of neighbours or a small domain, this will help us target specific variables. It may even turn out that these variables are exactly those for which testing the variable-elimination rules is less costly in computational resources. Our preliminary investigations we made in this direction seem to show that this is the case. Indeed, they establish that the eliminated variables often have a small domain or a small number of neighbours as shown in Figures 12 and 13. The two figures compare respectively the number of eliminated variables having a given domain size or a given degree (i.e. number of neighbours) to the corresponding number in the original instances. We focus our study on variables having a domain size or a degree at most 100. Above, the number of eliminated variables is negligible (about one percent in the best cases). We can observe that in our experiments the BT-degree property did not eliminate all variables with a singleton domain. This is simply explained by the instances for which the elimination process runs out of memory.

9.3 Impact on solving efficiency

This subsection is devoted to the impact of elimination rules on the solving efficiency. So, we consider the 1,337 instances for which at least one of the considered rules allows to eliminate some variables. For each of them, we apply MAC+RST+NG on the original instance and on the instance after (possibly) eliminating variables and we compare the observed runtime. In the latter case, the

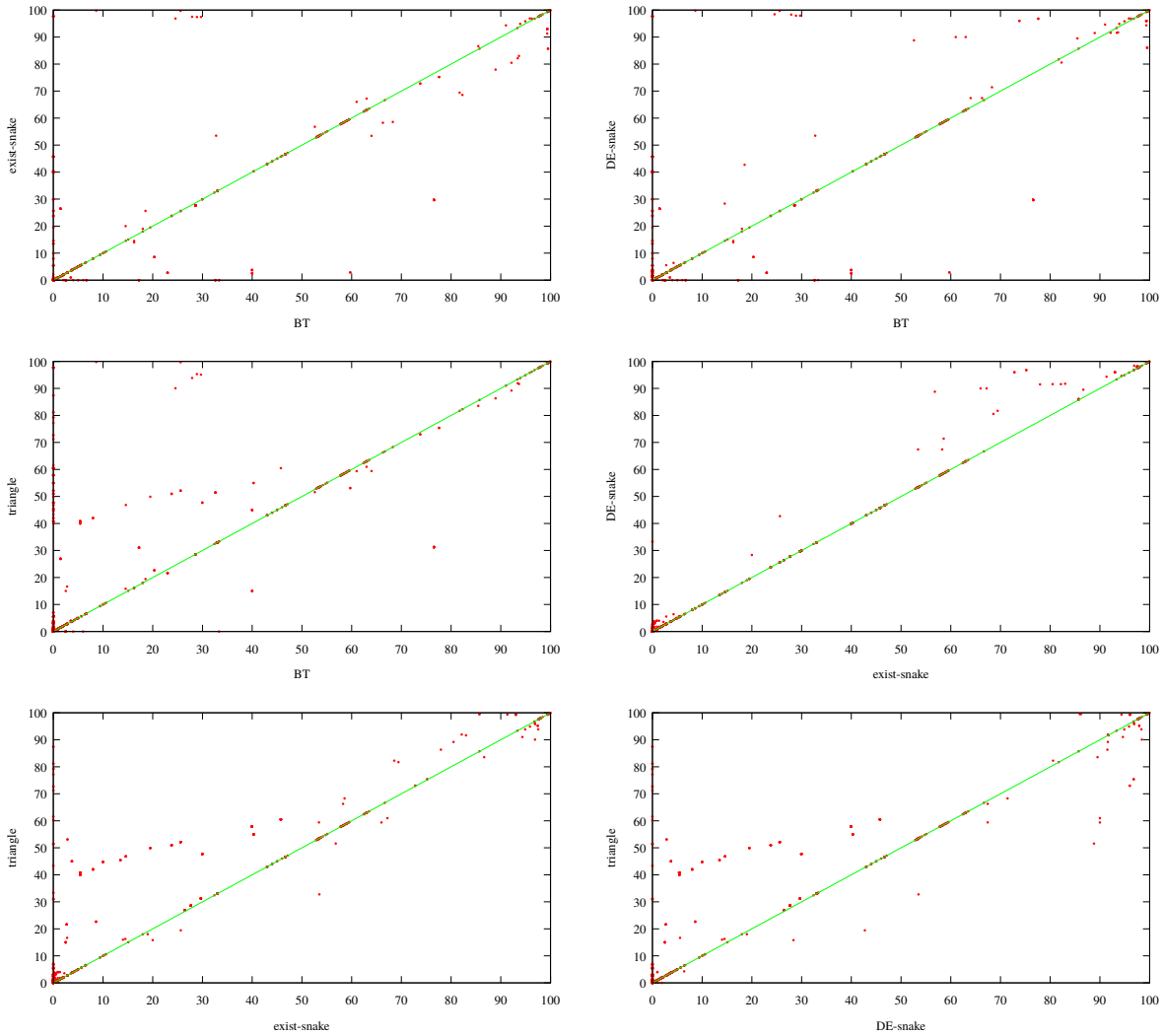


Figure 11: Comparison of the percentage of eliminated variables for each instance and for each pair of elimination rules.

runtime includes both the solving runtime and the variable-elimination phase runtime. Figure 14 gives the cumulative number of instances solved by MAC+RST+NG after eliminating some variables or by considering the original instances. The “step” which appears after 30 minutes in the BT curve is due to the fact that a large number of instances have just reached the time-out for the variable-elimination phase and are then solved fairly quickly. To compare fairly the different algorithms we have to observe the curves after this 30-minute mark. It appears that MAC+RST+NG solves more instances when the instances are preprocessed with any of the considered elimination rules. Moreover, the triangle property is again the most interesting elimination rule. Its use allows MAC+RST+NG to solve 1,008 instances while it only solves 996 and 998 instances when the instances are preprocessed respectively with the BT-degree property and \exists snake rule (1,000 instances for DE-snake rule). Without any preprocessing, MAC+RST+NG performs worst by solving only

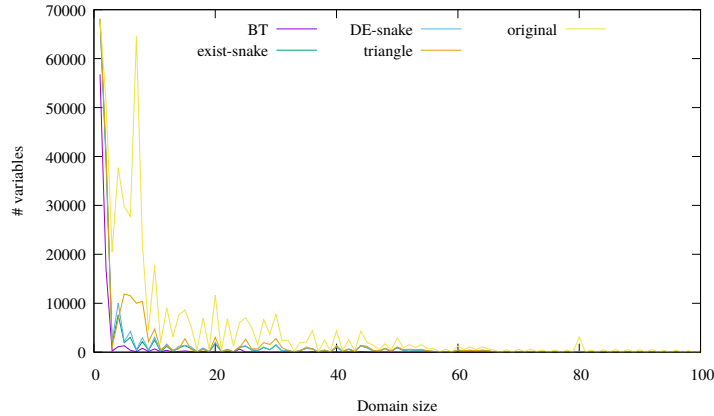


Figure 12: Number of eliminated variables having a domain of a given size for each elimination rule and number of such variables in the original instances.

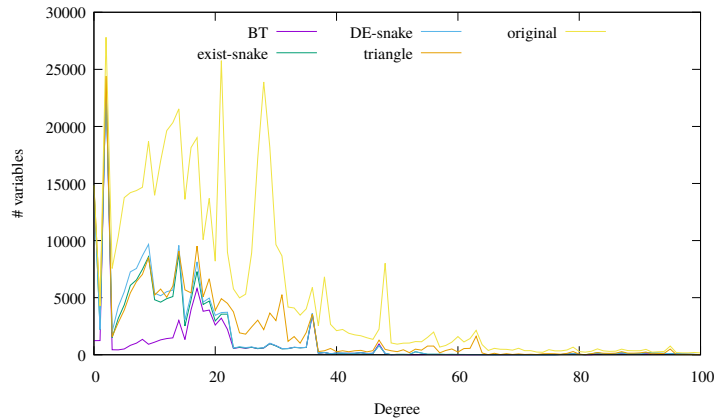


Figure 13: Number of eliminated variables having a given degree for each elimination rule and number of such variables in the original instances.

991 instances. If we compare more finely the runtime of MAC+RST+NG applied on the original instances and on the instances after eliminating some variables thanks to the triangle property (see Figure 15), we observe that, depending on the instance, eliminating variables may or may not improve solving efficiency. However, we can remark that there exist several instances which MAC+RST+NG solves after the elimination of some variables but not without and above all that the converse is false. So applying the triangle property for eliminating variables before solving makes sense.

Since the computational complexity of our variable-elimination rules is comparable with strong path consistency (SPC) (Lecoutre, 2009), it was natural to also test applying SPC. However, applying SPC in preprocessing allowed us to solve only 869 instances compared to 991 instances using MAC+RST+NG alone (without any variable elimination) and hence proved to be counter-productive. Applying SPC required the rewriting in extension of those relations that need to be modified, whereas our variable-elimination rules allow us to keep these relations in their original form.

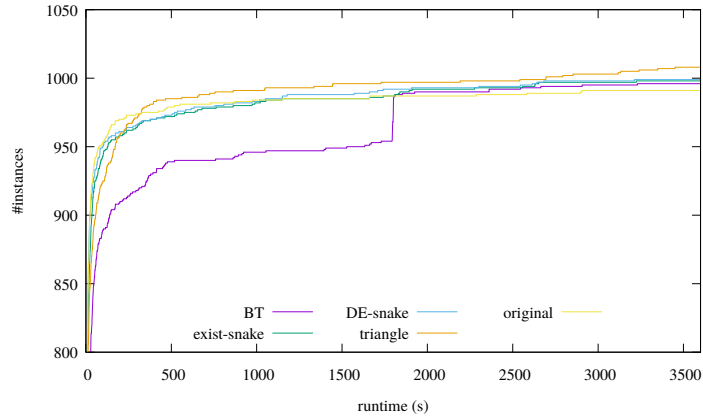


Figure 14: Cumulative number of instances solved by MAC+RST+NG after eliminating some variables or by considering the original instances with respect to elapsed time.

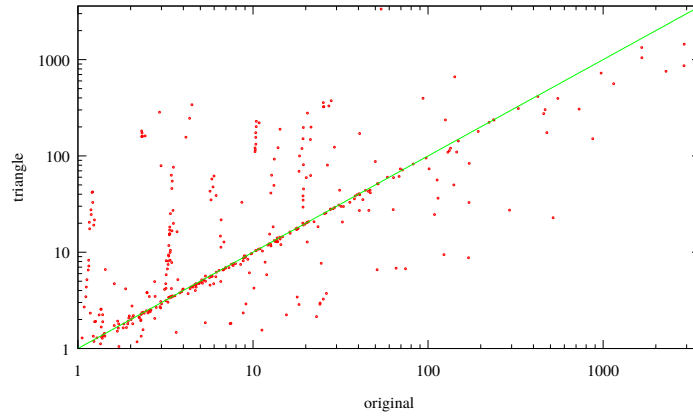


Figure 15: Comparison of the runtime of MAC+RST+NG on the original instances and on the same instances after eliminating variables thanks to the triangle property.

10. Variable-elimination rules and tractability

We investigate, in this section, the possibility of defining tractable classes based on our variable-elimination rules. As is the case for BTP (Cooper et al., 2010), the rules we have presented in this paper also define tractable classes that can be detected in polynomial time by successive elimination of variables.

We also study the confluence of our variable-elimination rules which allows us to show the tractability of maximising the number of eliminated variables.

Definition 14 *For a variable-elimination property P , we say that a binary CSP instance I satisfies P for the variable order $<$ if for each variable x_m , except for the first variable according to the order $<$, I satisfies the property P on x_m in the sub-instance of I restricted to the variables x_i such that $x_i \leq x_m$.*

Definition 15 We say that a property P of binary CSP instances is hereditary if for any instance I with more than one variable, I satisfies P implies that I_{-m} satisfies P , where I_{-m} is the instance obtained from I after elimination of the variable x_m .

Theorem 5 Let P be an hereditary sol-var-elim property which can be tested in polynomial time. The class of binary CSP instances I satisfying the property P (for a possibly unknown ordering of its variables) can be detected and solved in polynomial time.

Proof: Let I be a binary CSP instance on n variables. Suppose that I satisfies the hereditary variable-elimination property P for a variable ordering $<$. Then I satisfies property P on the last variable of the ordering $<$. We can therefore find a variable x_m on which I satisfies the property P by exhaustive search over all n variables. Note that there may be more than one variable which satisfies P . In this case, we make an arbitrary choice which variable to eliminate: the rest of the proof does require that x_m be the last variable according to the order $<$. Variable x_m is then eliminated to produce the sub-instance I_{-m} which has the same satisfiability as I . Since P is hereditary, the instance I_{-m} also satisfies P . By successive elimination of variables we can reduce I to an equivalent instance I' on a single variable in polynomial time. A single-variable instance being trivial to solve, and since P is a sol-var-elim property, we can construct a solution to I in polynomial time. ■

The following theorem is a direct consequence of Theorem 5 and the fact that the listed properties are hereditary.

Theorem 6 The class of binary CSP instances I satisfying any of the following properties (for a possibly unknown ordering of its variables) can be detected and solved in polynomial time:

1. the $\forall\exists$ broken k -dimensional polyhedron property (for any fixed $k \geq 2$),
2. the \exists snake property,
3. the DE-snake property,
4. $\forall\exists$ BTP,
5. the BT-degree property.

At first sight, it might appear that we would not have an equivalent result for the triangle property, since eliminating a variable x_j might destroy the triangle property on another variable x_i ($i \neq j$). Recall that, in Definition 5, a variable x_i can be eliminated by the triangle property only if there is a variable x_j which justifies this elimination, so the obvious question is whether x_i can still be eliminated after (its justifying variable) x_j has been eliminated. It turns out that the answer is yes, as we will now demonstrate.

Definition 16 In a binary CSP instance I , for distinct variable x_i, x_j , variable x_j justifies the elimination by the triangle property of variable x_i (which we denote $x_j \xrightarrow{I} x_i$) if for all $v_j \in \mathcal{D}(x_j)$, there exists $u_{ji}(v_j) \in \mathcal{D}(x_i)$ satisfying the following conditions:

$$\mathbf{C1}(i,j,v_j): (u_{ji}(v_j), v_j) \in R_{ij},$$

C2(i, j, v_j): $\forall x_k \in X \setminus \{x_i, x_j\}, \forall v_k \in \mathcal{D}(x_k), (v_j, v_k) \in R_{jk} \Rightarrow (u_{ji}(v_j), v_k) \in R_{ik}$.

Lemma 1 *If $x_p \xrightarrow{I} x_j$ and $x_j \xrightarrow{I} x_i$, where $p \neq i$, then $x_p \xrightarrow{I_{-j}} x_i$. In other words, if x_p justifies the elimination of x_j and x_j justifies the elimination of x_i by the triangle property, then x_p justifies the elimination of x_i in the instance I_{-j} which is the result of the elimination of x_j from I .*

Proof: Suppose that $x_p \xrightarrow{I} x_j$ and $x_j \xrightarrow{I} x_i$, where $p \neq i$. For $v_p \in \mathcal{D}(x_p)$, define

$$u_{pi}(v_p) := u_{ji}(u_{pj}(v_p)).$$

It suffices to show that $u_{pi}(v_p) \in \mathcal{D}(x_i)$ and that it satisfies the conditions of Definition 16 in I_{-j} , namely:

C1(i, p, v_p): $(u_{pi}(v_p), v_p) \in R_{ip}$,

C2(i, p, v_p): $\forall x_k \in X \setminus \{x_i, x_p, x_j\}, \forall v_k \in \mathcal{D}(x_k), (v_p, v_k) \in R_{pk} \Rightarrow (u_{pi}(v_p), v_k) \in R_{ik}$.

Since $x_p \xrightarrow{I} x_j$, we have $u_{pj}(v_p) \in \mathcal{D}(x_j)$ and then, since $x_j \xrightarrow{I} x_i$, we have $u_{pi}(v_p) \in \mathcal{D}(x_i)$ (with $u_{pi}(v_p) = u_{ji}(u_{pj}(v_p))$). Since $x_p \xrightarrow{I} x_j$, we also have from **C1**(j, p, v_p) that $(u_{pj}(v_p), v_p) \in R_{jp}$. Then, since $x_j \xrightarrow{I} x_i$, we have **C2**($i, j, u_{pj}(v_p)$) and, in particular for $k = p$ and $v_k = v_p \in \mathcal{D}(x_p)$:

$$(u_{pj}(v_p), v_p) \in R_{jp} \Rightarrow (u_{ji}(u_{pj}(v_p)), v_p) \in R_{ip}$$

Since $u_{pi}(v_p) = u_{ji}(u_{pj}(v_p))$, we can deduce that condition **C1**(i, p, v_p) holds.

Now consider any $x_k \in X \setminus \{x_i, x_p, x_j\}$ and any $v_k \in \mathcal{D}(x_k)$. Since $x_p \xrightarrow{I} x_j$, we have from **C2**(j, p, v_p) that

$$(v_p, v_k) \in R_{pk} \Rightarrow (u_{pj}(v_p), v_k) \in R_{jk}.$$

Since $x_j \xrightarrow{I} x_i$, we have from **C2**($i, j, u_{pj}(v_p)$):

$$(u_{pj}(v_p), v_k) \in R_{jk} \Rightarrow (u_{ji}(u_{pj}(v_p)), v_k) \in R_{ik}.$$

Hence, we have

$$(v_p, v_k) \in R_{pk} \Rightarrow (u_{ji}(u_{pj}(v_p)), v_k) \in R_{ik}.$$

Since $u_{pi}(v_p) = u_{ji}(u_{pj}(v_p))$, it follows that condition **C2**(i, p, v_p) holds, which completes the proof. ■

To complete our study of the definition of a tractable class based on the triangle property, we need to consider the one case not covered by Lemma 1, namely $p = i$. For this, we require the notions of isomorphic instances and neighbourhood substitutability of values. Indeed, when $x_i \xrightarrow{I} x_j$ and $x_j \xrightarrow{I} x_i$ we have to choose which of x_i or x_j we eliminate by the triangle property. We will show that modulo isomorphism (and provided we have applied neighbourhood substitution), the resulting instances are identical and hence we can make an arbitrary choice between x_i and x_j .

Definition 17 *Two binary CSP instances $I = \langle X^I, \mathcal{D}^I, R^I \rangle$, $J = \langle X^J, \mathcal{D}^J, R^J \rangle$ are isomorphic if there exist bijections $f : X^I \rightarrow X^J$ and $g_i : \mathcal{D}^I(x_i) \rightarrow \mathcal{D}^J(f(x_i))$ (for all $x_i \in X^I$) such that for all pairs of distinct variables $x_i, x_j \in X^I$, for all $v_i \in \mathcal{D}^I(x_i)$ and for all $v_j \in \mathcal{D}^I(x_j)$,*

$$(v_i, v_j) \in R_{ij}^I \Leftrightarrow (g_i(v_i), g_j(v_j)) \in R_{f(i)f(j)}^J$$

A neighbourhood-substitutable value can be eliminated from its domain without changing the satisfiability of the instance (Freuder, 1991; Freuder & Wallace, 2017).

Definition 18 *In a binary CSP instance $I = \langle X, \mathcal{D}, R \rangle$, $v_i \in \mathcal{D}(x_i)$ is neighbourhood substitutable by $v'_i \in \mathcal{D}(x_i) \setminus \{v_i\}$ if for all $x_j \in X \setminus \{x_i\}$, for all $v_j \in \mathcal{D}(x_j)$,*

$$(v_i, v_j) \in R_{ij} \Rightarrow (v'_i, v_j) \in R_{ij}$$

Two values v_i and v'_i are interchangeable if v_i is neighbourhood substitutable by v'_i and v'_i is neighbourhood substitutable by v_i .

Given a binary CSP instance I , it is known that the result of eliminating neighbourhood-substitutable values until convergence (i.e. no more eliminations are possible) is unique up to isomorphism (Cooper, 1997). We assume that there is a program NS which performs value eliminations by neighbourhood substitutability until convergence, and we denote by $NS(I)$ the result of applying NS to I .

Lemma 2 *If $x_i \xrightarrow{I} x_j$ and $x_j \xrightarrow{I} x_i$ then $NS(I_{-i})$ and $NS(I_{-j})$ are isomorphic.*

Proof: It follows from the definition of $x_i \xrightarrow{I} x_j$ that for all $v_i \in \mathcal{D}(x_i)$, we have $u_{ij}(v_i) \in \mathcal{D}(x_j)$ such that for all $x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$,

$$(v_i, v_k) \in R_{ik} \Rightarrow (u_{ij}(v_i), v_k) \in R_{jk} \quad (1)$$

Since $x_j \xrightarrow{I} x_i$ we then have $u_{ji}(u_{ij}(v_i)) \in \mathcal{D}(x_i)$ such that for all $x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$,

$$(u_{ij}(v_i), v_k) \in R_{jk} \Rightarrow (u_{ji}(u_{ij}(v_i)), v_k) \in R_{ik} \quad (2)$$

For each $x_i \in X$, define the function $F_i : \mathcal{D}(x_i) \rightarrow \mathcal{D}(x_i)$ by $F_i(v_i) = u_{ji}(u_{ij}(v_i))$ and consider the sequence

$$v_i, F(v_i), F(F(v_i)), \dots, F^r(v_i), \dots$$

Since $\mathcal{D}(x_i)$ is finite, this sequence must cycle at some point. Let $r \in \mathbb{N}$ be the first value for which $F^r(v_i) = F^s(v_i)$ for some $s > r$. Thus, the above sequence has a cycle of length $s - r$ starting at $F^r(v_i)$.

From Equations 1 and 2, we know that for all $x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$,

$$(v_i, v_k) \in R_{ik} \Rightarrow (F(v_i), v_k) \in R_{ik}$$

By a simple inductive argument, we can deduce that for any $t \in \mathbb{N}$,

$$(F^t(v_i), v_k) \in R_{ik} \Rightarrow (F^r(v_i), v_k) \in R_{ik} \quad (3)$$

and hence each $F^t(v_i) \neq F^r(v_i)$ is neighbourhood substitutable by $F^r(v_i)$ in I_{-j} . Thus, modulo isomorphism, we can assume that all values in the sequence $v_i, F(v_i), F^2(v_i), \dots$ have been eliminated from $\mathcal{D}(x_i)$ in $NS(I_{-j})$ by neighbourhood substitution except for $F^r(v_i)$ (Cooper, 1997). By a similar argument, modulo isomorphism, we can assume that in $NS(I_{-i})$ all values in the sequence $u_{ij}(v_i), u_{ij}(F(v_i)), u_{ij}(F^2(v_i)), \dots$ have been eliminated by neighbourhood substitutability from

$\mathcal{D}(x_j)$ except for $u_{ij}(F^r(v_i))$. Furthermore, combining Equations 1, 2 and 3, we can deduce that for all $x_k \in X \setminus \{x_i, x_j\}$, $\forall v_k \in \mathcal{D}(x_k)$,

$$(F^r(v_i), v_k) \in R_{ik} \Leftrightarrow (u_{ij}(F^r(v_i)), v_k) \in R_{jk}$$

Thus, for each value in $\mathcal{D}(x_i)$ in $NS(I_{-j})$, there is a corresponding value in $\mathcal{D}(x_j)$ in $NS(I_{-i})$ which has the same compatibilities with all values for all other variables (and vice versa). Furthermore, no two values in $\mathcal{D}(x_i)$ (respectively, $\mathcal{D}(x_j)$) can have the same compatibilities with all values for all other variables, otherwise they would be interchangeable in $NS(I_{-j})$ (respectively, $NS(I_{-i})$) which would contradict the definition of neighbourhood substitution. It follows that $NS(I_{-i})$ and $NS(I_{-j})$ are isomorphic. ■

We require one final lemma.

Lemma 3 *If $x_i \xrightarrow{I} x_j$ then $x_i \xrightarrow{NS(I)} x_j$.*

Proof: Suppose that $x_i \xrightarrow{I} x_j$. We can see from Definition 5 that no eliminations of values from $\mathcal{D}(x_i)$ or $\mathcal{D}(x_k)$ (for $x_k \in X \setminus \{x_i, x_j\}$) can possibly invalidate the elimination of x_j by the triangle property. Suppose that $v_j \in \mathcal{D}(x_j)$ is eliminated in $NS(I)$ since it is neighbourhood substitutable by $v'_j \in \mathcal{D}(x_i)$. If $v_j = u_{ij}(v_i)$ for some $v_i \in \mathcal{D}(x_i)$, it suffices to set the value of $u_{ij}(v_i)$ to v'_j instead of v_j . It is easy to see that neighbourhood substitutability guarantees that conditions **C1**(j, i, v_i) and **C2**(j, i, v_i) in Definition 16 hold with this new value of $u_{ij}(v_i)$. ■

We can now prove that the triangle property defines a tractable class which is detectable in polynomial time.

Theorem 7 *The class of binary CSP instances I satisfying the triangle property (for a possibly unknown ordering of its variables) can be detected and solved in polynomial time.*

Proof: Suppose that there exists a variable order $<$ for which I satisfies the triangle property. Let x_m be the last variable according to this (unknown) order. We can find in $O(\text{end}^3)$, using the algorithm in Appendix C the set T of variables which could be eliminated from I by (a first pass of) the triangle property. We know that $T \neq \emptyset$ since $x_m \in T$. We do not know which variable in T is the last variable according to the order $<$, so we eliminate some arbitrary variable $x_i \in T$ from I . We then perform neighbourhood-substitution eliminations until convergence to obtain $I' = NS(I_{-i})$. By Lemma 1, x_m can be eliminated by the triangle property from I' except possibly in the case that $x_i \xrightarrow{I} x_m$ and $x_m \xrightarrow{I} x_i$. But, in this latter case, by Lemma 2, I' is isomorphic to $NS(I_{-m})$ and so it is as if we had eliminated x_m instead of x_i . Lemma 3 tells us that that eliminating values by neighbourhood substitutability does not destroy the fact that an instance satisfies the triangle property. We can deduce that the instance I' satisfies the triangle property and hence, by an easy inductive argument, that we will reduce the instance to a single-variable instance by successive eliminations of $n - 1$ variables. The theorem follows from Theorem 2. ■

When not all variables can be eliminated, we are interested in maximising the number of eliminated variables. As pointed out in the proof of Theorem 5, the elimination of a variable by a hereditary rule cannot be invalidated by the elimination of another variable. The following theorem is an immediate consequence of this and the fact that the listed properties are hereditary.

Theorem 8 *Maximising the number of variables that can be eliminated by any of the following rules can be achieved in polynomial time: the $\forall\exists$ broken k -dimensional polyhedron property (for any fixed $k \geq 2$), the \exists snake property, the DE-snake property, $\forall\exists$ BTP and the BT-degree property.*

We saw in the proof of Theorem 7 that the elimination of a variable x_m by the triangle property can only be invalidated by the elimination of another variable x_i by the triangle property if $NS(I_{-i})$ is isomorphic to $NS(I_{-m})$. It follows that the triangle property is confluent modulo isomorphism, provided neighbourhood substitution is applied after every variable elimination. We thus have the following theorem.

Theorem 9 *Maximising the number of variable eliminations by combining the triangle property and neighbourhood substitution can be achieved in polynomial time.*

11. Discussion and conclusion

In this paper we have given novel satisfiability-conserving variable-elimination rules for binary CSPs, two of which (namely DE-snake and BT-degree) strengthen previously-published rules. In each case, if the instance is satisfiable, then a solution to the original instance can be recovered in low-order polynomial time from a solution to the reduced instance. We have given optimised algorithms for applying each rule until convergence. The DE-snake rule can be applied until convergence in $O(ed^3)$ time, whereas the corresponding time complexity for the triangle rule and the BT-degree rule is $O(end^3)$. However, it should be pointed out that the DE-snake rule inherits the disadvantage of the \exists snake rule that the number of solutions may actually increase after elimination of a variable (Cohen et al., 2015): for example, it allows us to eliminate the central variable in the two-colouring of a star graph which increases the number of solutions from 2 to 2^{n-1} .

Extensive experimental trials have confirmed that because of relatively high time complexity of each of the variable-elimination rules, they may only be tested exhaustively during preprocessing. Applying them in preprocessing allowed us to solve more benchmark instances than without, with the triangle rule allowing us to eliminate more variables and hence solve more instances than the other rules. From a practical viewpoint, it would be interesting to understand how to better target the instances or the variables for which the proposed variable elimination rules are likely to be profitable. As a first step in this direction we have seen that most variables eliminated by our rules have small domain size and/or small degree. Future work is required to determine whether versions of our rules targetting only certain variables may be a practical possibility during search.

We have, in particular, generalised the notion of broken triangle to broken polyhedron, which may be of independent theoretical interest. The broken polyhedra property may lead to other possible theoretical advances (such as value-merging (Cooper et al., 2016a), value-elimination (Cohen et al., 2015), and generalisations to the general-arity CSP (Cooper et al., 2016a; El Mouelhi, 2017) or the Quantified CSP (Gao, Yin, & Zhou, 2011)), as was the case with the broken-triangle property (Cooper et al., 2010).

We have also shown that each of the variable-elimination rules allows us to define a novel hybrid tractable class by successive elimination of almost all variables. For each rule, this elimination order can be found in polynomial time, which we found surprising in the case of the triangle property.

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Appendix A. Algorithm for variable elimination by the \exists snake property

In this and the following appendices, we assume that we have implemented a set data structure in such a way that we can perform the following operations in $O(1)$ time: set membership, addition/deletion of an element and testing whether the set is empty. This can be achieved using a boolean table together with a counter of the number of elements in the set, since in each case the set is a subset of a fixed set, such as X the variables of the instance.

Below we give an algorithm for eliminating variables by the \exists snake property until convergence. It uses the following data structures :

- S_{ELIM} is the set of variables to be eliminated.
- For $v_j, v'_j \in \mathcal{D}(x_j)$, $\text{vars}_-^+(j, v_j, v'_j) = \{k \mid \exists v_k \in \mathcal{D}(x_k) \text{ with } (v_j, v_k) \in R_{jk} \wedge (v'_j, v_k) \notin R_{jk}\}$.
- For distinct $i, j \in \{1, \dots, n\}$ such that x_i constrains x_j and $v_i \in \mathcal{D}(x_i)$, $\text{countPairs}(i, v_i, j)$ is the number of pairs of values $v_j, v'_j \in \mathcal{D}(x_j)$ such that for some $k \neq i, j$ and some $v_k \in \mathcal{D}(x_k)$, the snake pattern (as shown in Figure 1) occurs on $\langle x_i, v_i \rangle, \langle x_j, v_j \rangle, \langle x_j, v'_j \rangle, \langle x_k, v_k \rangle$. We calculate $\text{countPairs}(i, v_i, j)$ by noting that it is the number of pairs of values $v_j, v'_j \in \mathcal{D}(x_j)$ such that $(v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji} \wedge \text{vars}_-^+(j, v_j, v'_j) \setminus \{i\} \neq \emptyset$.
- For $v_i \in \mathcal{D}(x_i)$, $\text{badVars}(i, v_i)$ is the set of $j \neq i$ such that $\text{countPairs}(i, v_i, j) \neq 0$. If $\text{badVars}(i, v_i) = \emptyset$, then variable x_i is added to S_{ELIM} .

The algorithm first initialises the above data structures, then performs eliminations from X , the set of variables (which is initially $\{x_1, \dots, x_n\}$). When performing an elimination, the data structures are updated which may lead to more variable eliminations. Eliminations propagate until convergence (i.e. until no more eliminations are possible). Updating the data structures, when a variable x_k is eliminated, means deleting k from each $\text{vars}_-^+(j, v_j, v'_j)$ and each $\text{badVars}(i, v_i)$. When deleting k from $\text{vars}_-^+(j, v_j, v'_j)$, the value of $\text{countPairs}(i, v_i, j)$ needs to be decremented only in the case that $((v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji})$ and $\text{vars}_-^+(j, v_j, v'_j) \setminus \{i\}$ becomes empty for the first time (i.e. $\text{vars}_-^+(j, v_j, v'_j)$ becomes $\{i\}$ or it becomes empty and the k being deleted from it is not i).

*** Initialisation ***

$S_{ELIM} := \emptyset$;

for $x_j \in X$:

 for $v_j \in \mathcal{D}(x_j)$:

 for $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$:

 for $x_k \in X \setminus \{x_j\}$ such that x_k is constrained by x_j :

 if $\exists v_k \in \mathcal{D}(x_k)$ such that $(v_j, v_k) \in R_{jk} \wedge (v'_j, v_k) \notin R_{jk}$

 then add k to $\text{vars}_-^+(j, v_j, v'_j)$;

```

for  $x_i \in X$  :
  for  $v_i \in \mathcal{D}(x_i)$  : .....(1)
    badVars( $i, v_i$ ) :=  $\emptyset$  ;
    for  $x_j \in X \setminus \{x_i\}$  such that  $x_j$  is constrained by  $x_i$  :
      countPairs( $i, v_i, j$ ) := 0 ;
      for  $v_j \in \mathcal{D}(x_j)$  :
        for  $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$  :
          if  $(v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji} \wedge \text{vars}_-^+(j, v_j, v'_j) \setminus \{i\} \neq \emptyset$ 
          then countPairs( $i, v_i, j$ ) := countPairs( $i, v_i, j$ ) + 1 ;
            badVars( $i, v_i$ ) := badVars( $i, v_i$ )  $\cup$   $\{j\}$  ;
        if badVars( $i, v_i$ ) =  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ; exit loop (1) ;

```

*** Elimination and propagation ***

```

while  $S_{ELIM} \neq \emptyset$  :
  delete some  $x_k$  from  $S_{ELIM}$  ;  $X := X \setminus \{x_k\}$  ;
  for  $x_j \in X$  such that  $x_j$  is constrained by  $x_k$  :
    for  $v_j \in \mathcal{D}(x_j)$  :
      for  $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$  :
        delete  $k$  from  $\text{vars}_-^+(j, v_j, v'_j)$  ;
        if  $\text{vars}_-^+(j, v_j, v'_j)$  becomes a singleton  $\{i\}$  after deletion of  $k$  .....(2)
        then for  $v_i \in \mathcal{D}(x_i)$  :
          if  $((v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji})$ 
          then countPairs( $i, v_i, j$ ) := countPairs( $i, v_i, j$ ) - 1 ;
            if countPairs( $i, v_i, j$ ) = 0
            then badVars( $i, v_i$ ) := badVars( $i, v_i$ )  $\setminus$   $\{j\}$  ;
              if badVars( $i, v_i$ ) =  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;
          if  $\text{vars}_-^+(j, v_j, v'_j)$  becomes  $\emptyset$  after deletion of  $k$  .....(3)
        then for  $x_i \in X \setminus \{x_j, x_k\}$  such that  $x_i$  is constrained by  $x_j$  :
          for  $v_i \in \mathcal{D}(x_i)$  :
            if  $((v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji})$ 
            then countPairs( $i, v_i, j$ ) := countPairs( $i, v_i, j$ ) - 1 ;
              if countPairs( $i, v_i, j$ ) = 0
              then badVars( $i, v_i$ ) := badVars( $i, v_i$ )  $\setminus$   $\{j\}$  ;
                if badVars( $i, v_i$ ) =  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;
          for  $i \in X$  such that  $x_i$  is constrained by  $x_k$  :
            for  $v_i \in \mathcal{D}(x_i)$  : .....(4)
              badVars( $i, v_i$ ) := badVars( $i, v_i$ )  $\setminus$   $\{k\}$  ;
              if badVars( $i, v_i$ ) =  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ; exit loop (4) ;

```

This algorithm requires $O(ed^3)$ time and $O(ed^2)$ space. To see the $O(ed^3)$ time bound, observe that each of the tests (2) and (3) can only be True once for each triple (j, v_j, v'_j) . The data structure $\text{vars}_-^+(j, v_j, v'_j)$ requires $O(ed^2)$ space.

Appendix B. Algorithm for variable elimination by the DE-snake property

Below we give an algorithm for eliminating variables by the DE-snake property until convergence. It uses the following data structures :

- S_{ELIM} is the set of variables to be eliminated.
- For $v_j, v'_j \in \mathcal{D}(x_j)$, $\text{vars}_-^+(j, v_j, v'_j) = \{k \mid \exists v_k \in \mathcal{D}(x_k) \text{ with } (v_j, v_k) \in R_{jk} \wedge (v'_j, v_k) \notin R_{jk}\}$.
- For $v_i \in \mathcal{D}(x_i)$, $\text{badAssts}(i, v_i)$ is the set of assignments $\langle j, v_j \rangle$ such that $(v_i, v_j) \notin R_{ij}$ and $\nexists v'_j \in \mathcal{D}(x_j)$ such that $\text{vars}_-^+(j, v_j, v'_j) \subseteq \{i\}$ and $(v_i, v'_j) \in R_{ij}$. If $\text{badAssts}(i, v_i) = \emptyset$, then variable x_i is added to S_{ELIM} since it can be eliminated by the DE-snake rule.

The algorithm first initialises the above data structures, then performs eliminations from X , the set of variables (which is initially $\{x_1, \dots, x_n\}$). When performing an elimination, the data structures are updated which may lead to more variable eliminations. Eliminations propagate until convergence (i.e. until no more eliminations are possible). Updating the data structures, when a variable x_k is eliminated, means deleting k from each $\text{vars}_-^+(j, v_j, v'_j)$ and deleting each assignment to x_k from each $\text{badAssts}(i, v_i)$. When deleting k from $\text{vars}_-^+(j, v_j, v'_j)$, $\text{badAssts}(i, v_i)$ needs to be updated for each variable x_i for which $\text{vars}_-^+(j, v_j, v'_j) \setminus \{i\}$ becomes empty for the first time (i.e. $\text{vars}_-^+(j, v_j, v'_j)$ becomes $\{i\}$ or it becomes empty and the k being deleted from it is not i).

*** Initialisation ***

$S_{ELIM} := \emptyset$;

for $x_j \in X$:

 for $v_j \in \mathcal{D}(x_j)$:

 for $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$:

 for $x_k \in X \setminus \{x_j\}$ such that x_k is constrained by x_j :

 if $\exists v_k \in \mathcal{D}(x_k)$ such that $(v_j, v_k) \in R_{jk} \wedge (v'_j, v_k) \notin R_{jk}$

 then add k to $\text{vars}_-^+(j, v_j, v'_j)$;

for $x_i \in X$:

 for $v_i \in \mathcal{D}(x_i)$:

$\text{badAssts}(i, v_i) := \emptyset$

 for x_j such that x_j is constrained by x_i :

 for $v_j \in \mathcal{D}(x_j)$:

 if $(v_i, v_j) \notin R_{ij}$

 then for $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$:

 if $(v_i, v'_j) \in R_{ij}$ and $\text{vars}_-^+(j, v_j, v'_j) \setminus \{i\} = \emptyset$

 then add $\langle j, v_j \rangle$ to $\text{badAssts}(i, v_i)$;

 if $\text{badAssts}(i, v_i) = \emptyset$ then add x_i to S_{ELIM} ;

*** Elimination and propagation ***

```

while  $S_{ELIM} \neq \emptyset$  :
  delete some  $x_k$  from  $S_{ELIM}$  ;  $X := X \setminus \{x_k\}$  ;
  for  $x_j \in X$  such that  $x_j$  is constrained by  $x_k$  :
    for  $v_j \in \mathcal{D}(x_j)$  :
      for  $v'_j \in \mathcal{D}(x_j) \setminus \{v_j\}$  :
        delete  $k$  from  $\text{vars}^+(j, v_j, v'_j)$  ;
        if  $\text{vars}^+(j, v_j, v'_j)$  becomes a singleton  $\{i\}$  after deletion of  $k$  : .....(1)
        then for  $v_i \in \mathcal{D}(x_i)$  :
          if  $((v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji})$ 
          then  $\text{badAssts}(i, v_i) := \text{badAssts}(i, v_i) \setminus \{(j, v_j)\}$  ;
          if  $\text{badAssts}(i, v_i)$  becomes  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;
        if  $\text{vars}^+(j, v_j, v'_j)$  becomes  $\emptyset$  after deletion of  $k$  : .....(2)
        then for  $x_i \in X \setminus \{x_j, x_k\}$  such that  $x_i$  is constrained by  $x_j$  :
          for  $v_i \in \mathcal{D}(x_i)$  :
            if  $((v'_j, v_i) \in R_{ji} \wedge (v_j, v_i) \notin R_{ji})$ 
            then  $\text{badAssts}(i, v_i) := \text{badAssts}(i, v_i) \setminus \{(j, v_j)\}$  ;
            if  $\text{badAssts}(i, v_i)$  becomes  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;
    for  $i \in X$  such that  $x_i$  is constrained by  $x_k$  :
      for  $v_i \in \mathcal{D}(x_i)$  :
        for  $v_k \in \mathcal{D}(x_k)$  :
           $\text{badAssts}(i, v_i) := \text{badVars}(i, v_i) \setminus \{(k, v_k)\}$  ;
          if  $\text{badVars}(i, v_i)$  becomes  $\emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;

```

This algorithm requires $O(ed^3)$ time and $O(ed^2)$ space. To see the $O(ed^3)$ time bound, observe that each of the tests (1) and (2) can only be True once for each triple (j, v_j, v'_j) . The data structures $\text{vars}^+(j, v_j, v'_j)$ and $\text{badAssts}(i, v_i)$ both require $O(ed^2)$ space.

Appendix C. Algorithm for variable elimination by the triangle property

Below we give an algorithm for eliminating variables by the triangle property until convergence. It uses the following data structures :

- S_{ELIM} is the set of variables to be eliminated.
- For $v_j \in \mathcal{D}(x_j)$, $v_i \in \mathcal{D}(x_i)$ such that $(v_j, v_i) \in R_{ji}$, $\text{badVars}(j, v_j, i, v_i)$ is the set of $k \neq i, j$ such that there exists $v_k \in \mathcal{D}(x_k)$ with $(v_j, v_k) \in R_{jk}$ and $(v_i, v_k) \notin R_{ik}$.
- $\text{supported}(j, v_j, i) = \text{True}$ if there exists $v_i \in \mathcal{D}(x_i)$ such that $\text{badVars}(j, v_j, i, v_i) = \emptyset$.
- $\text{count}(j, i)$ is the number of values $v_j \in \mathcal{D}(x_j)$ such that $\text{supported}(j, v_j, i) = \text{False}$. If $\text{count}(j, i) = 0$, for some $j \neq i$ such that $x_j \in X \setminus S_{ELIM}$ then we can eliminate x_i by the triangle property.

The algorithm first initialises the above data structures, then performs eliminations. When performing an elimination, these data structures are updated which may lead to further variable eliminations. The only delicate point in the algorithm is that once a variable is due to be eliminated it cannot be used to justify the elimination of another variable; hence the test $x_j \notin S_{ELIM}$ in lines (1) and (3).


```

*** Initialisation ***
 $S_{ELIM} := \emptyset$ ;
for  $x_i \in X$  :
  for  $x_j \in X \setminus \{x_i\}$  such that  $x_j \notin S_{ELIM}$  : .....(1)
    count( $j, i$ ) := 0 ;
    for  $v_j \in \mathcal{D}(x_j)$ 
      supported( $j, v_j, i$ ) := False ;
      for  $v_i \in \mathcal{D}(x_i)$  such that  $(v_j, v_i) \in R_{ji}$  : .....(2)
        badVars( $j, v_j, i, v_i$ ) :=  $\emptyset$  ;
        for  $x_k \in X \setminus \{x_i, x_j\}$  such that  $x_k$  is constrained by  $x_i$  :
          if  $\exists v_k \in \mathcal{D}(x_k)$  such that  $(v_j, v_k) \in R_{jk}$  and  $(v_i, v_k) \notin R_{ik}$ 
            then add  $k$  to badVars( $j, v_j, i, v_i$ ) ;
          if badVars( $j, v_j, i, v_i$ ) =  $\emptyset$ 
            then supported( $j, v_j, i$ ) := True ; exit loop (2) ;
          if supported( $j, v_j, i$ ) = False then count( $j, i$ ) := count( $j, i$ ) + 1 ;
        if count( $j, i$ ) = 0
          then add  $x_i$  to  $S_{ELIM}$  ; exit loop (1) ;

```

```

*** Elimination and propagation ***
while  $S_{ELIM} \neq \emptyset$  :
  delete some  $x_k$  from  $S_{ELIM}$  ;  $X := X \setminus \{x_k\}$  ;
  for  $x_i \in X$  such that  $x_i$  is constrained by  $x_k$  :
    for  $x_j \in X \setminus \{x_i\}$  such that  $x_j \notin S_{ELIM}$  : .....(3)
      for  $v_j \in \mathcal{D}(x_j)$  such that supported( $j, v_j, i$ ) = False :
        for  $v_i \in \mathcal{D}(x_i)$  : .....(4)
          if  $k \in \text{badVars}(j, v_j, i, v_i)$ 
            then delete  $k$  from badVars( $j, v_j, i, v_i$ )
            if badVars( $j, v_j, i, v_i$ ) =  $\emptyset$ 
              then supported( $j, v_j, i$ ) := True ;
              count( $j, i$ ) := count( $j, i$ ) - 1 ;
              if count( $j, i$ ) = 0
                then add  $x_i$  to  $S_{ELIM}$  ; exit loop (3) ;
            exit loop (4) ;

```

This algorithm requires $O(\text{end}^3)$ time and $O(\text{end}^2)$ space. Curiously, the propagation phase requires less time than initialisation phase, since it requires only $O(\text{end}^2)$ time. This is because we no longer need to look at individual values in the propagation phase.

Appendix D. Algorithm for variable elimination by the BT-degree property

Below we give an algorithm for eliminating variables by the BT-degree property until convergence. It uses the following data structures :

- S_{ELIM} is the set of variables to be eliminated.

- $N_-^+(i, v_i, j, v_j, m) = |\{v_m \in \mathcal{D}(x_m) \mid (v_i, v_m) \in R_{im} \wedge (v_j, v_m) \notin R_{jm}\}|$, is the number of values $v_m \in \mathcal{D}(x_m)$ which are linked by a positive edge to $\langle x_i, v_i \rangle$ and by a negative edge to $\langle x_j, v_j \rangle$.
- For a pair of variables x_i, x_m linked by a constraint, for $v_i \in \mathcal{D}(x_i)$ and $v_m \in \mathcal{D}(x_m)$, $\text{BTvars}(i, v_i, m, v_m)$ is the set of $j \neq i, m$ such that there is a broken triangle (v_m, v_i, v_j, v'_m) or (v_m, v_j, v_i, v'_m) for some $v_j \in \mathcal{D}(x_j)$ and $v'_m \in \mathcal{D}(x_m)$.
- $\text{BTdegree}(i, v_i, m, v_m)$ is the cardinality of the set $\text{BTvars}(i, v_i, m, v_m)$.
- $\text{3safe}(i, v_i, j, v_j, m)$ is True if $(v_i, v_j) \in R_{ij}$ is 3-safe on x_m .
- $M_-^+(i, v_i, j, v_j, m)$ is similar to $N_-^+(i, v_i, j, v_j, m)$, except that the negative edge $v_j v_m$ must also have BT degree greater than 1.
- $\text{badBases}(m)$ is the set of (i, v_i, j, v_j) such that i, j, m are distinct, $v_i \in \mathcal{D}(x_i)$, $v_j \in \mathcal{D}(x_j)$ and $\nexists v'_m \in \mathcal{D}(x_m)$ satisfying the conditions in the definition of the BT-degree property (namely, $(v_i, v'_m) \in R_{im}$, $(v_j, v'_m) \in R_{jm}$ and either (v_i, v_j) is 3-safe on x_m or (v_i, v'_m) has BT degree zero or (v_j, v'_m) has BT degree zero). Thus, if $\text{badBases}(m)$ is the empty set, then x_m can be eliminated by the BT-degree property.

The algorithm first initialises the above data structures, then performs eliminations. Again, the set of variables X is initially $\{x_1, \dots, x_n\}$. The data structure 3safe is calculated using the fact that $(v_i, v_j) \in R_{ij}$ is 3-safe on x_m if and only if $M_-^+(i, v_i, j, v_j, m) = 0$ or $M_-^+(j, v_j, i, v_i, m) = 0$. When performing an elimination, all data structures are updated which may provoke further variable eliminations.

*** Initialisation ***

$S_{ELIM} := \emptyset$;

for $x_m \in X$:

 for $x_i \in X \setminus \{x_m\}$ such that x_i is constrained by x_m :

 for $x_j \in X \setminus \{x_i, x_m\}$ such that x_j is constrained by x_m :

 for $v_i \in \mathcal{D}(x_i)$:

 for $v_j \in \mathcal{D}(x_j)$ such that $(v_i, v_j) \in R_{ij}$:

$N_-^+(i, v_i, j, v_j, m) := |\{v_m \in \mathcal{D}(x_m) \mid (v_i, v_m) \in R_{im} \wedge (v_j, v_m) \notin R_{jm}\}|$;

for $x_m \in X$:

 for $x_i \in X \setminus \{x_m\}$ such that x_i is constrained by x_m :

 for $v_i \in \mathcal{D}(x_i)$:

 for $v_m \in \mathcal{D}(x_m)$:

$\text{BTvars}(i, v_i, m, v_m) := \{j \neq i, m \mid \exists v_j \in \mathcal{D}(x_j)$ such that $(v_i, v_j) \in R_{ij}$ and

$((v_i, v_m) \notin R_{im} \wedge (v_j, v_m) \in R_{jm} \wedge N_-^+(i, v_i, j, v_j, m) \neq 0)$

 or $((v_i, v_m) \in R_{im} \wedge (v_j, v_m) \notin R_{jm} \wedge N_-^+(j, v_j, i, v_i, m) \neq 0)\}$;

$\text{BTdegree}(i, v_i, m, v_m) := |\text{BTvars}(i, v_i, m, v_m)|$;

for $x_m \in X$:

$\text{badBases}(m) := \emptyset$;

 for $x_i \in X \setminus \{x_m\}$ such that x_i is constrained by x_m :

 for $x_j \in X \setminus \{x_i, x_m\}$ such that x_j is constrained by x_m :

for $v_i \in \mathcal{D}(x_i)$:
 for $v_j \in \mathcal{D}(x_j)$ such that $(v_i, v_j) \in R_{ij}$:(1)
 $M_-^+(i, v_i, j, v_j, m) := |\{v_m \in \mathcal{D}(x_m) \mid (v_i, v_m) \in R_{im} \wedge (v_j, v_m) \notin R_{jm} \wedge \text{BTdegree}(j, v_j, m, v_m) > 1\}|$;
 $3\text{safe}(i, v_i, j, v_j, m) := M_-^+(i, v_i, j, v_j, m) = 0$ or $M_-^+(j, v_j, i, v_i, m) = 0$;
 if $\nexists v'_m \in \mathcal{D}(x_m)$ such that $(v_i, v'_m) \in R_{im} \wedge (v_j, v'_m) \in R_{jm} \wedge$
 $(3\text{safe}(i, v_i, j, v_j, m) \vee \text{BTdegree}(i, v_i, m, v'_m) = 0 \vee \text{BTdegree}(j, v_j, m, v'_m) = 0)$
 then add (i, v_i, j, v_j) to $\text{badBases}(m)$; exit loop (1) ;
 if $\text{badBases}(m) = \emptyset$ then add x_m to S_{ELIM} ;

*** Elimination and propagation ***

while $S_{ELIM} \neq \emptyset$:
 delete some x_j from S_{ELIM} ; $X := X \setminus \{x_j\}$;
 for $x_m \in X$ such that x_m is constrained by x_j :
 for $x_i \in X \setminus \{x_m\}$ such that x_i is constrained by x_m :
 for $v_i \in \mathcal{D}(x_i)$:
 for $v_m \in \mathcal{D}(x_m)$:
 if $j \in \text{BTvars}(i, v_i, m, v_m)$
 then delete j from $\text{BTvars}(i, v_i, m, v_m)$;
 $\text{BTdegree}(i, v_i, m, v_m) := \text{BTdegree}(i, v_i, m, v_m) - 1$;
 if $\text{BTdegree}(i, v_i, m, v_m) = 1$ (2)
 then for $x_k \in X \setminus \{x_i, x_m\}$ such that x_k is constrained by x_m :
 for $v_k \in \mathcal{D}(x_k)$ such that $(v_i, v_k) \in R_{ik}$:
 if $(v_k, v_m) \in R_{km} \wedge (v_i, v_m) \notin R_{im}$
 then $M_-^+(k, v_k, i, v_i, m) := M_-^+(k, v_k, i, v_i, m) - 1$;
 if $(M_-^+(i, v_i, k, v_k, m) = 0 \vee M_-^+(k, v_k, i, v_i, m) = 0)$ becomes True
 then $3\text{safe}(i, v_i, k, v_k, m) := \text{True}$;
 for $v'_m \in \mathcal{D}(x_m)$ such that $(v_i, v'_m) \in R_{im} \wedge (v_k, v'_m) \in R_{km}$:
 $\text{badBases}(m) := \text{badBases}(m) \setminus \{(i, v_i, k, v_k)\}$;
 if $\text{badBases}(m) = \emptyset$ then add x_m to S_{ELIM} ;
 if $\text{BTdegree}(i, v_i, m, v_m) = 0$ (3)
 then for $x_k \in X \setminus \{x_i, x_m\}$ such that x_k is constrained by x_m :
 for $v_k \in \mathcal{D}(x_k)$ such that $(v_i, v_k) \in R_{ik}$:
 if $(v_i, v_m) \in R_{im}$ and $(v_k, v_m) \in R_{km}$
 then $\text{badBases}(m) := \text{badBases}(m) \setminus \{(i, v_i, k, v_k)\}$;
 if $\text{badBases}(m) = \emptyset$ then add x_m to S_{ELIM} ;

This algorithm requires $O(\text{end}^3)$ time and $O(\text{end}^2 \log d)$ space. To prove the $O(\text{end}^3)$ time bound we have to use the fact that each of the tests (2) and (3) can only become True once for each quadruple (i, v_i, m, v_m) . The data structure badBases requires $O(\text{end}^2)$ space and the data structures N_-^+ and M_-^+ require $O(\text{end}^2 \log d)$ space, the $\log d$ factor being due to the fact that they store integers up to d .

Appendix E. Algorithm for variable elimination by $\forall\exists\text{BTP}$

Below we give an algorithm for eliminating variables by $\forall\exists\text{BTP}$ until convergence. It uses the following data structures :

- S_{ELIM} is the set of variables to be eliminated.
- $\text{badVars}(i)$ is the set of $j \neq i$ such that there is some value $v_j \in \mathcal{D}(x_j)$ such that for all values $v_i \in \mathcal{D}(x_i)$ there is a broken triangle (v_i, v_j, v_k, v'_i) on x_i for some v'_i, v_k, k such that $v'_i \in \mathcal{D}(x_i)$ and $v_k \in \mathcal{D}(x_k)$.
- $\text{countBadVals}(j, i)$ is the number of values $v_j \in \mathcal{D}(x_j)$ such that for all values $v_i \in \mathcal{D}(x_i)$ there is a broken triangle (v_i, v_j, v_k, v'_i) on x_i for some v'_i, v_k, k such that $v'_i \in \mathcal{D}(x_i)$ and $v_k \in \mathcal{D}(x_k)$.
- $\text{support}(j, v_j, i)$ is the set of all values $v_i \in \mathcal{D}(x_i)$ such that there is no broken triangle (v_i, v_j, v_k, v'_i) on x_i for any v'_i, v_k, k with $v'_i \in \mathcal{D}(x_i)$ and $v_k \in \mathcal{D}(x_k)$.
- $L_{BT}(j, v_j, i, v_i)$ is the set of k such that there is a broken triangle (v_i, v_j, v_k, v'_i) on x_i for some v'_i, v_k, k such that $v'_i \in \mathcal{D}(x_i)$ and $v_k \in \mathcal{D}(x_k)$.

The algorithm first initialises the above data structures, then performs eliminations. When performing an elimination, these data structures are updated which may provoke further variable eliminations.

*** Initialisation ***

$S_{ELIM} := \emptyset$;

for $x_i \in X$:

$\text{badVars}(i) := \emptyset$;

 for $x_j \in X \setminus \{x_i\}$ such that x_j is constrained by x_i :

$\text{countBadVals}(j, i) := 0$;

 for $v_j \in \mathcal{D}(x_j)$:

$\text{support}(j, v_j, i) := \emptyset$; $L := \emptyset$;

 for $x_k \in X \setminus \{x_i, x_j\}$ such that x_k is constrained by x_i :

 for $v_k \in \mathcal{D}(x_k)$ such that $(v_j, v_k) \in R_{jk}$:

 if $\exists v'_i \in \mathcal{D}(x_i)$ such that $(v_j, v'_i) \notin R_{ji}$ and $(v'_i, v_k) \in R_{ik}$

 then add $\langle k, v_k \rangle$ to L ;

 for $v_i \in \mathcal{D}(x_i)$ such that $(v_j, v_i) \in R_{ji}$:

$L_{BT}(j, v_j, i, v_i) := \emptyset$;

 for $\langle k, v_k \rangle \in L$:

 if $(v_i, v_k) \notin R_{ik}$ then $L_{BT}(j, v_j, i, v_i) := L_{BT}(j, v_j, i, v_i) \cup \{k\}$;

 if $L_{BT}(j, v_j, i, v_i) = \emptyset$ then add v_i to $\text{support}(j, v_j, i)$;

 if $\text{support}(j, v_j, i) = \emptyset$

 then $\text{countBadVals}(j, i) := \text{countBadVals}(j, i) + 1$; $\text{badVars}(i) := \text{badVars}(i) \cup \{j\}$;

 if $\text{badVars}(i) = \emptyset$ then add x_i to S_{ELIM} ;

*** Elimination and propagation ***

```

while  $S_{ELIM} \neq \emptyset$  :
  delete some variable  $x_k$  from  $S_{ELIM}$  ;  $X := X \setminus \{x_k\}$  ;
  for  $x_i \in X$  such that  $x_i$  is constrained by  $x_k$  :
    if  $k \in \text{badVars}(i)$  then delete  $k$  from  $\text{badVars}(i)$  ;
    for  $x_j \in X \setminus \{x_i, x_k\}$  such that  $x_j$  is constrained by  $x_i$  :
      for  $v_j \in \mathcal{D}(x_j)$  :
        for  $v_i \in \mathcal{D}(x_i)$  such that  $(v_j, v_i) \in R_{ji}$  :
          if  $k \in L_{BT}(j, v_j, i, v_i)$ 
            then delete  $k$  from  $L_{BT}(j, v_j, i, v_i)$  ;
          if  $L_{BT}(j, v_j, i, v_i) = \emptyset$ 
            then add  $v_i$  to  $\text{support}(j, v_j, i)$  ;
            if  $|\text{support}(j, v_j, i)| = 1$ 
              then  $\text{countBadVals}(j, i) := \text{countBadVals}(j, i) - 1$  ;
              if  $\text{countBadVals}(j, i) = 0$  then delete  $j$  from  $\text{badVars}(i)$  ;
        if  $\text{badVars}(i) = \emptyset$  then add  $x_i$  to  $S_{ELIM}$  ;

```

This algorithm requires $O(\text{end}^3)$ time and $O(\text{end}^2)$ space. The data structure L_{BT} requires $O(\text{end}^2)$ space.

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